

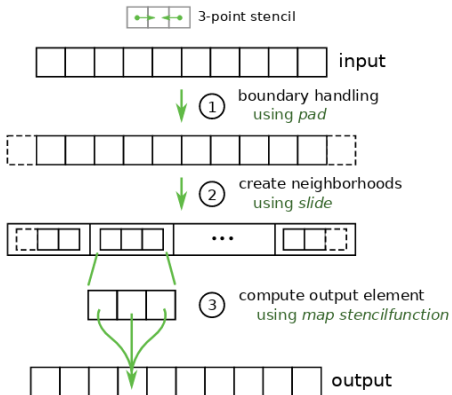
Lift Tutorial: View System

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Views

A Simple Expression

```
val highLevel = fun(
  ArrayType(Float, N), input =>
  Map(Reduce(add, 0.0f)) o
  Slide(3,1) o
  Pad(1,1,clamp) $ input )
```



Data Layout Primitives

```
val highLevel = fun(
  ArrayType(Float, N), input =>
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```

Observations:

- *Pad* and *Slide* only modify the data layout
 - How to avoid unnecessary temporary result?
- *Slide* increases the dimension of our one-dimensional input array
 - How to generate accesses to multi-dimensional arrays with a flat representation in memory?

Data Layout Primitives

```
val lowLevel = fun(
  ArrayType(Float, N), input =>
    MapGlb(MapSeq(toGlobal(id)) o ReduceSeq(add, 0.0f)) o
    Slide(3,1) o
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3-Point Stencil Code

```
float add(float x, float y) { return x + y; }
float id(float x) { return x; }
kernel void KERNEL(const global float *restrict IN, global float *OUT, int N) {
    float acc;
    // Map
    for (int globalID = get_global_id(0); (globalID < N);
         globalID = (globalID + get_global_size(0))) {
        acc = 0.0f;
        // Reduce
        for (int i = 0; i < 3; i++) {

            acc = add(acc, IN[???]);
        }
        OUT[globalID] = id(acc);
    }
}
```

Introducing Views

```
val lowLevel = fun(
  ArrayType(Float, N), input =>
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    Slide(3,1) o
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```

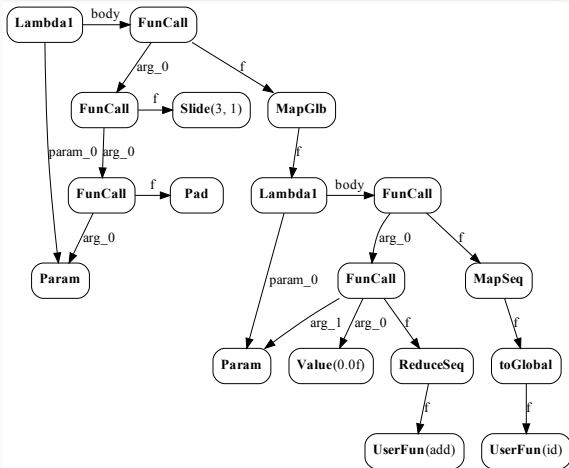
Views:

- **Construct** a representation of the effects of data layout functions
- **Consume** the views to generate correct array indices

```

val lowLevel = fun(
  ArrayType(Float, N), input =>
    MapGlb(MapSeq(toGlobal(id)) o ReduceSeq(add, 0.0f)) o
    Slide(3,1) o
    Pad(1,1,clamp) $ input )

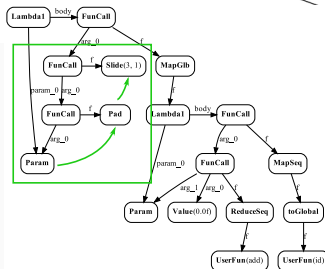
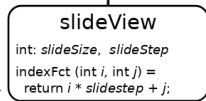
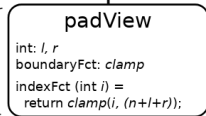
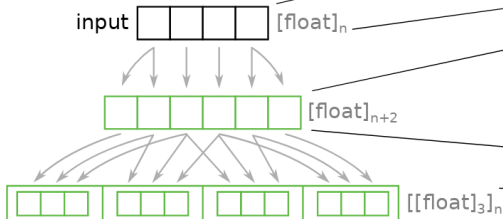
```



View Construction: Follow the Dataflow

$(\text{mapGlobal}(\text{reduceSeq} (+) 0) \circ \text{slide } 3 \ 1 \circ \text{pad } 1 \ 1 \ \text{clamp}) \text{ input}$

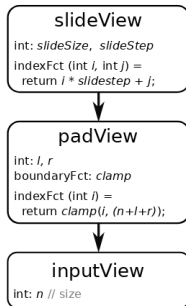
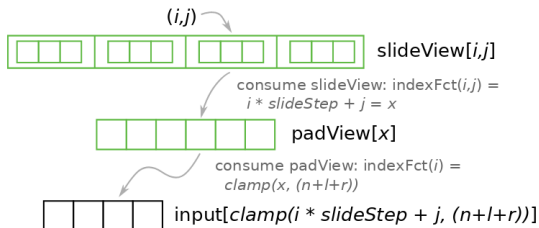
CONSTRUCTION ↓



View Consumption: Generate Flat Accesses

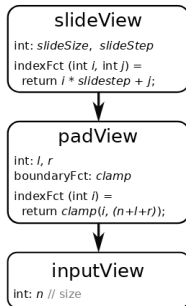
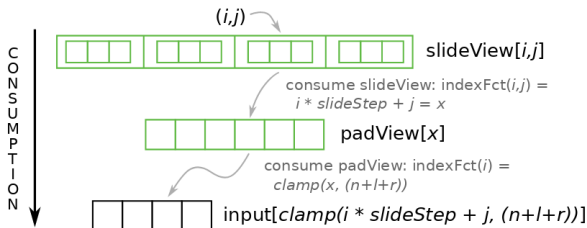
$(\text{mapGlobal}(\text{reduceSeq } (+) 0) \circ \text{slide } 3 \ 1 \circ \text{pad } 1 \ 1 \ \text{clamp}) \text{ input}$

CONSUMPTION
↓



View Consumption: Generate Flat Accesses

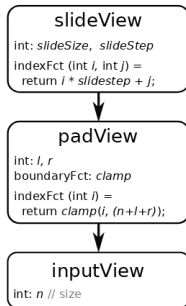
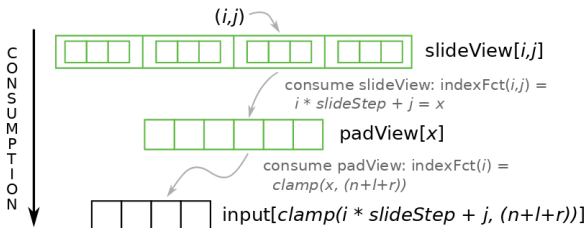
$(\text{mapGlobal}(\text{reduceSeq } (+) 0) \circ \text{slide } 3 \ 1 \circ \text{pad } 1 \ 1 \ \text{clamp}) \text{ input}$



Demo: emitView for ViewSlide

View Consumption: Generate Flat Accesses

$(\text{mapGlobal}(\text{reduceSeq } (+) 0) \circ \text{slide } 3 \ 1 \circ \text{pad } 1 \ 1 \ \text{clamp}) \text{ input}$



Demo: emitView for ViewSlide

```
for(int i = 0; i < 3; i++) {
    acc = add(acc, IN[((( $-1 + \text{globalID} + i \geq 0$ ) ?
        (( $-1 + \text{globalID} + i < N$ ) ?
            ( $-1 + \text{globalID} + i$ ) : ( $-1 + N$ )) : 0))]
}
```

View Consumption: Generate Flat Accesses

```
// Reduce unrolled  
acc = add(acc, IN[(((-1 + globalID + 0) ≥ 0) ? (((-1 + globalID + 0) < N) ? (-1 + globalID + 0) : (-1 + N)) : 0)]);  
acc = add(acc, IN[(((-1 + globalID + 1) ≥ 0) ? (((-1 + globalID + 1) < N) ? (-1 + globalID + 1) : (-1 + N)) : 0)]);  
acc = add(acc, IN[(((-1 + globalID + 2) ≥ 0) ? (((-1 + globalID + 2) < N) ? (-1 + globalID + 2) : (-1 + N)) : 0)]);
```

Are all these operations necessary or can we do better?

Arithmetic Expression Simplification

Lift comes with a powerful ArithExpr Library.

- performs simple arithmetic simplifications ($1 + 1 = 2$)
- keeps track of range information for variables
 - e.g., $0 \leq \text{globalID} < N$
- handles arithmetic operations including integer division and modulo
 - e.g., $((2M + 1) \bmod M) = 1 \bmod M$

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- performs simple arithmetic simplifications ($1 + 1 = 2$)
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 - e.g., $((2M + 1) \bmod M) = 1 \bmod M$

Demo:

1. Library
2. ArithExpr Type Hierarchy
3. Cst, Var (including ranges), ?, Mod
4. SimplifySum
5. Examples

Generated Indices Revisited

Question: Can we simplify the second array access?

```
// Reduce unrolled  
acc = add(acc, IN[((( $-1 + \text{globalID} + 0$ )  $\geq 0$ ) ? ((( $-1 + \text{globalID} + 0$ ) < N) ? ( $-1 + \text{globalID} + 0$ ) : ( $-1 + N$ )) : 0)]]);  
acc = add(acc, IN[((( $-1 + \text{globalID} + 1$ )  $\geq 0$ ) ? ((( $-1 + \text{globalID} + 1$ ) < N) ? ( $-1 + \text{globalID} + 1$ ) : ( $-1 + N$ )) : 0)]]);  
acc = add(acc, IN[((( $-1 + \text{globalID} + 2$ )  $\geq 0$ ) ? ((( $-1 + \text{globalID} + 2$ ) < N) ? ( $-1 + \text{globalID} + 2$ ) : ( $-1 + N$ )) : 0)]]);
```

Arithmetic Simplification

```
IN[ // predicate
    (((-1 + globalID + 1) >= 0) ?

    // true
    (((-1 + globalID + 1) < N) ?
    (-1 + globalID + 1) : (-1 + N)) :

    // false
    0)
];
```

Arithmetic Simplification

```
IN[ // predicate
    (((-1 + globalID + 1) >= 0) ?

    // true
    (((-1 + globalID + 1) < N) ?
    (-1 + globalID + 1) : (-1 + N)) :

    // false
    0)
];
```

Additions with constants cancel out

Arithmetic Simplification

```
IN[ // predicate
    ((globalID >= 0) ?

    // true
    ((globalID < N) ?
    globalID : (-1 + N)) :

    // false
    0)
];
```

Arithmetic Simplification

```
IN[ // predicate
    ((globalID >= 0) ?

    // true
    ((globalID < N) ?
    globalID : (-1 + N)) :

    // false
    0)
];
```

Predicate is always true (requires range information about the variable)

Arithmetic Simplification

```
//      predicate           true           false  
IN[(globalID < N) ? globalID : (-1 + N)];
```

Arithmetic Simplification

```
//for (int globalID = get_global_id(0);  
//     globalID < N; ...  
  
//     predicate           true           false  
IN[(globalID < N) ? globalID : (-1 + N)];
```

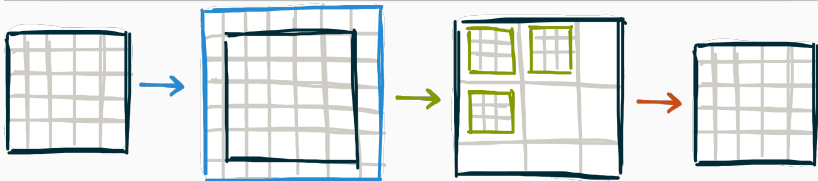
Arithmetic Simplification

```
IN[globalID];
```


Conclusion

```
val lowLevelExpression = fun(
  ArrayType(ArrayType(Float, M), N), input =>

  MapGlb(1)(MapGlb(0)(
    MapSeq(toGlobal(id)) o
      ReduceSeqUnroll(add, 0.0f) o Join()
  )) o Slide2D(3,1) o
  Pad2D(1,1,clamp) $ input)
```



Conclusion

```
kernel void KERNEL(const global float* restrict IN, global float* OUT, int M, int N){
    float acc;
    for (int y = get_global_id(1); (y < N); y = (y + get_global_size(1))) {
        for (int x = get_global_id(0); (x < M); x = (x + get_global_size(0))) {
            acc = 0.0f;
            // NW
            acc += IN[(M * ((-1 + y) ≥ 0) ? (-1 + y) : 0) + ((-1 + x) ≥ 0) ? (-1 + x) : 0)];
            // N
            acc += IN[(x + (M * ((-1 + y) ≥ 0) ? (-1 + y) : 0))];
            // NE
            acc += IN[(M * ((-1 + y) ≥ 0) ? (-1 + y) : 0) + ((1 + x) < M) ? (1 + x) : (-1 + M) )]);
            // W
            acc += IN[(M * y) + ((-1 + x) ≥ 0) ? (-1 + x) : 0)];
            // C
            acc += IN[(x + (M * y))];
            // E
            acc += IN[(M * y) + ((1 + x) < M) ? (1 + x) : (-1 + M) )]);
            // SW
            acc += IN[(M * ((1 + y) < N) ? (1 + y) : (-1 + N) ) + ((-1 + x) ≥ 0) ? (-1 + x) : 0)];
            // S
            acc += IN[(x + (M * ((1 + y) < N) ? (1 + y) : (-1 + N) ))]);
            // SE
            acc += IN[(M * ((1 + y) < N) ? (1 + y) : (-1 + N) ) + ((1 + x) < M) ? (1 + x) : (-1 + M) )]);
            // write back result
            OUT[(x + (M * y))] = acc;
        }
    }
}
```