

Lift

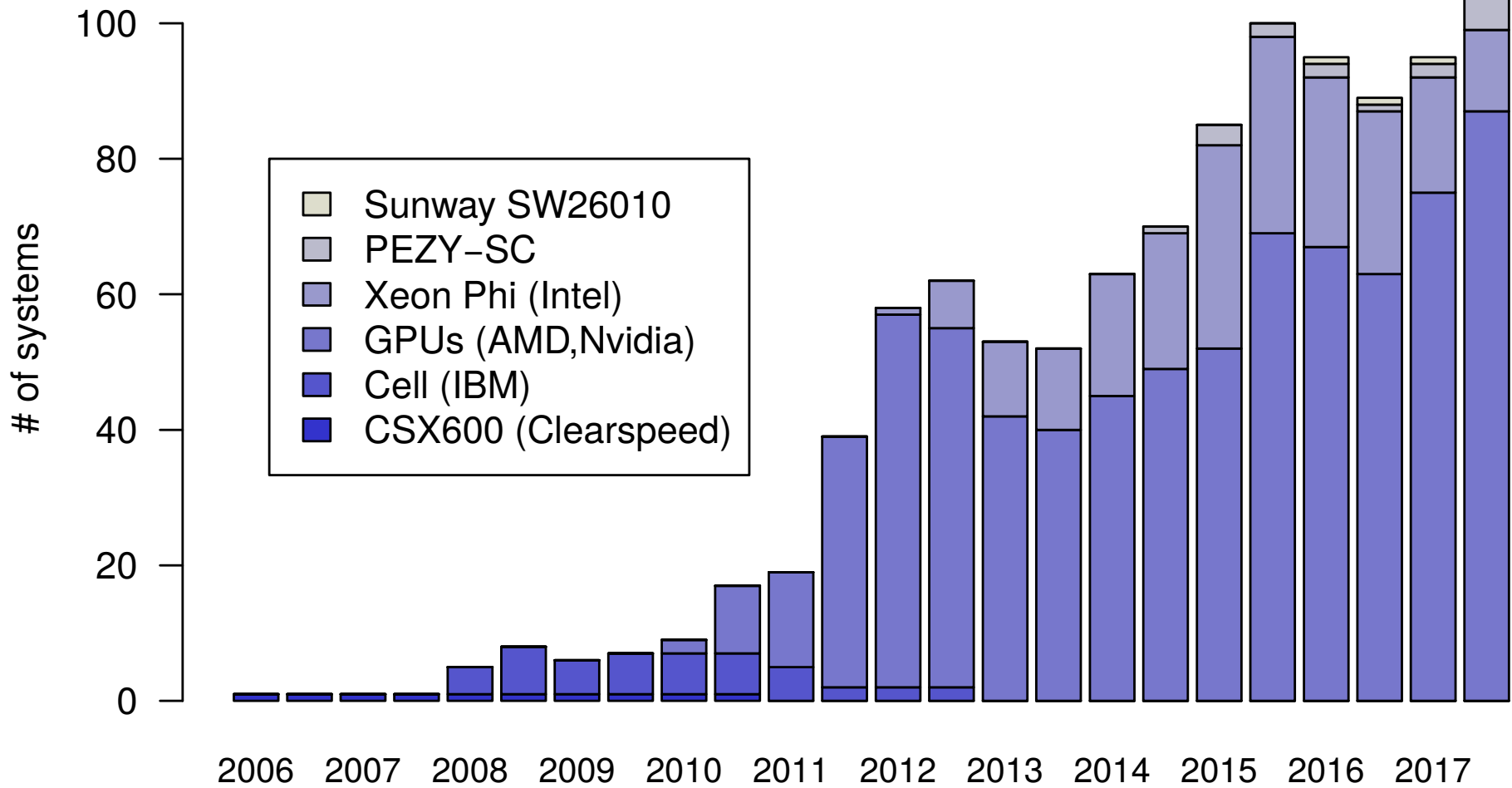
Performance Portable Code Generation on Parallel Accelerators



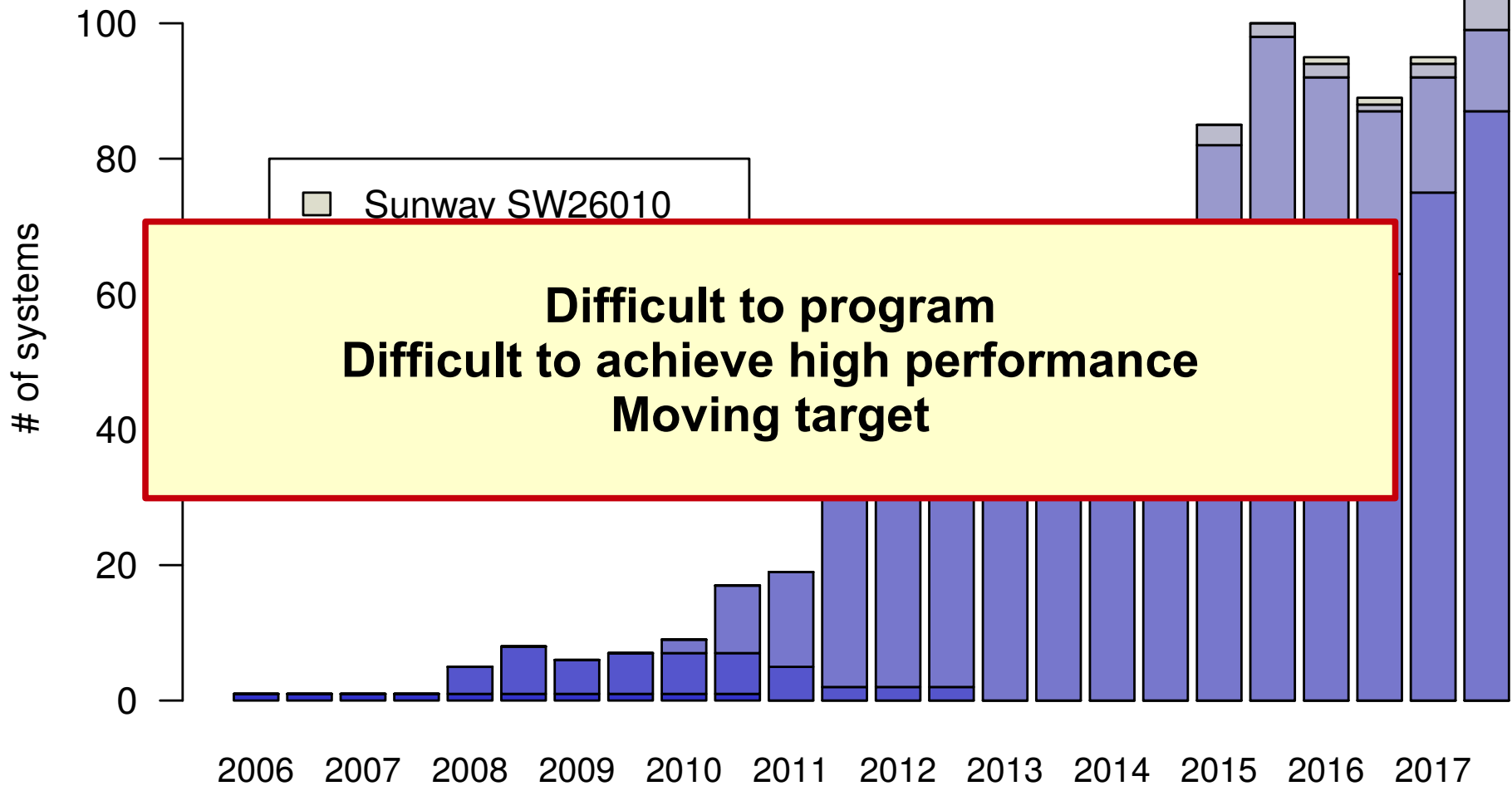
ISPASS tutorial

2 Apr. 2018

Top 500 with parallel accelerators



Top 500 with parallel accelerators



How to sum an array?

How to sum an array?

```
float acc = 0;  
for (int i=0; i<N; i++)  
    acc += input[i];
```

How to really sum an array:

```
kernel void sum(global float* g_in, global float* g_out,
               unsigned int n, local volatile float* l_data) {

    unsigned int tid = get_local_id(0);
    unsigned int i   = get_group_id(0) * 256 + get_local_id(0);

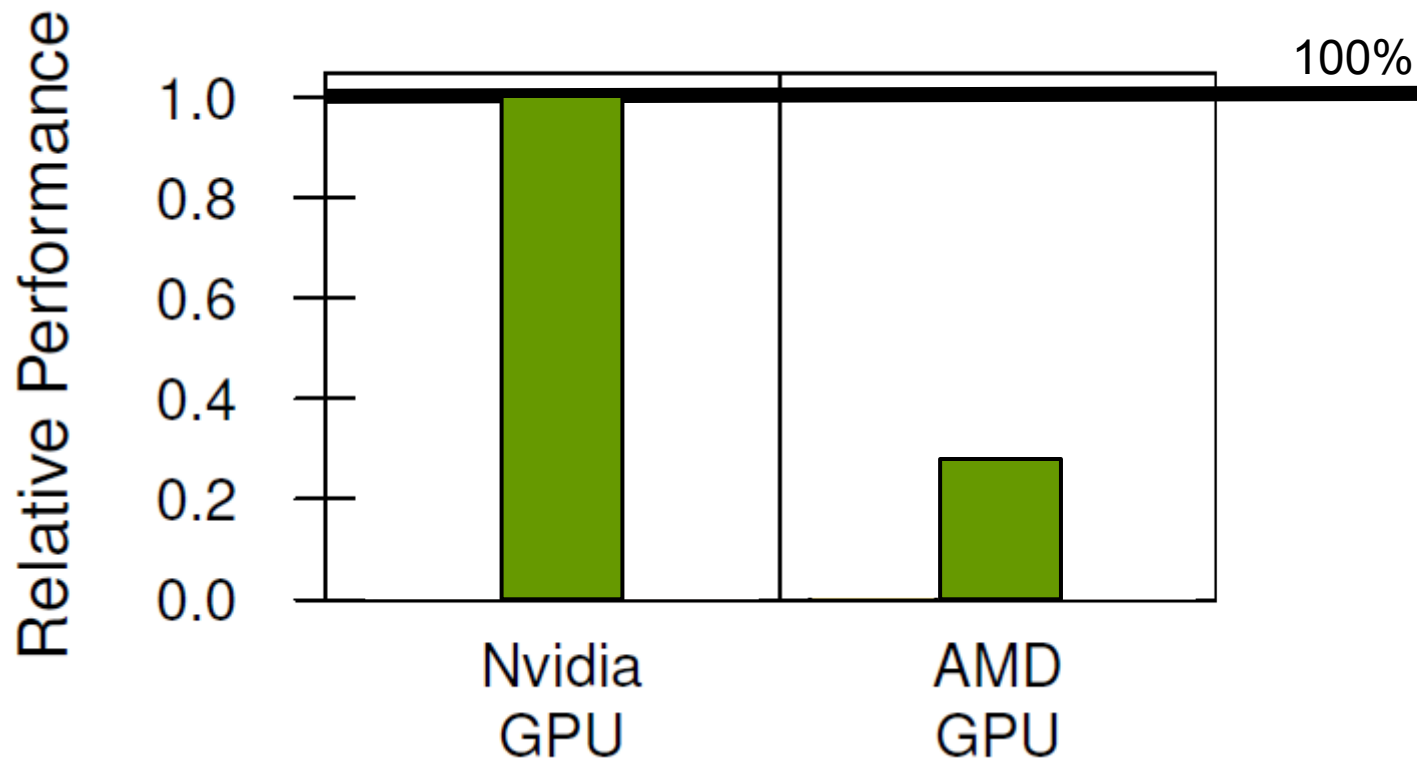
    l_data[tid] = 0;
    while (i < n) {
        l_data[tid] += g_in[i];
        i += 256 * get_num_groups(0);
    }
    barrier(CLK_LOCAL_MEM_FENCE);

    if (tid < 128)
        l_data[tid] += l_data[tid+128];
    barrier(CLK_LOCAL_MEM_FENCE);
    if (tid < 64)
        l_data[tid] += l_data[tid+ 64];
    barrier(CLK_LOCAL_MEM_FENCE)

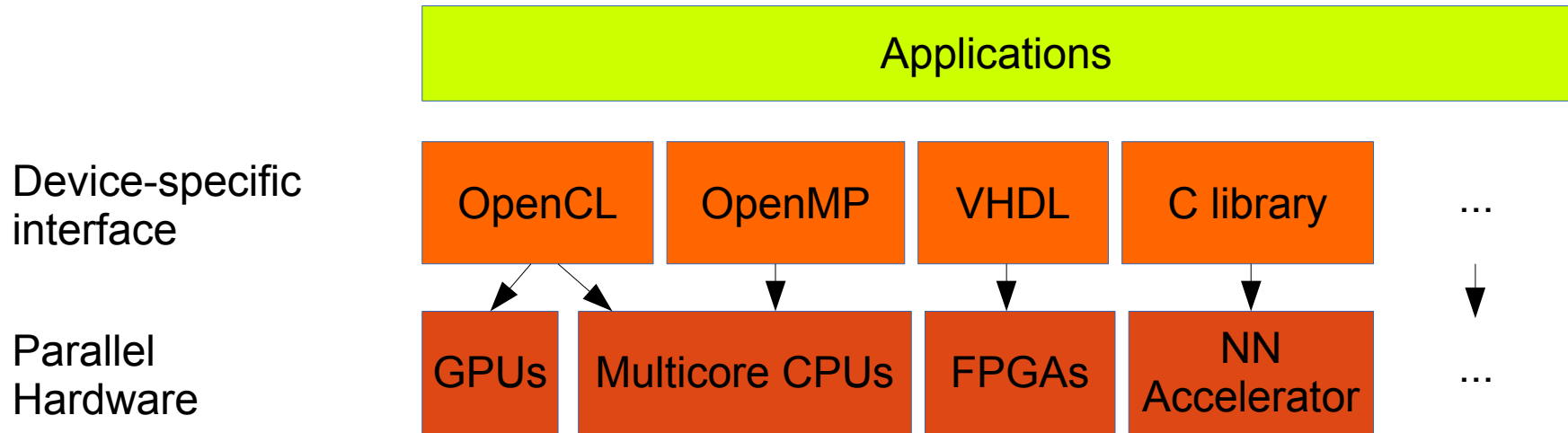
    if (tid < 32) {
        l_data[tid] += l_data[tid+32]; l_data[tid] += l_data[tid+16];
        l_data[tid] += l_data[tid+ 8]; l_data[tid] += l_data[tid+ 4];
        l_data[tid] += l_data[tid+ 2]; l_data[tid] += l_data[tid+ 1];
    }
    if (tid == 0)
        g_out[get_group_id(0)] = l_data[0];
}
```



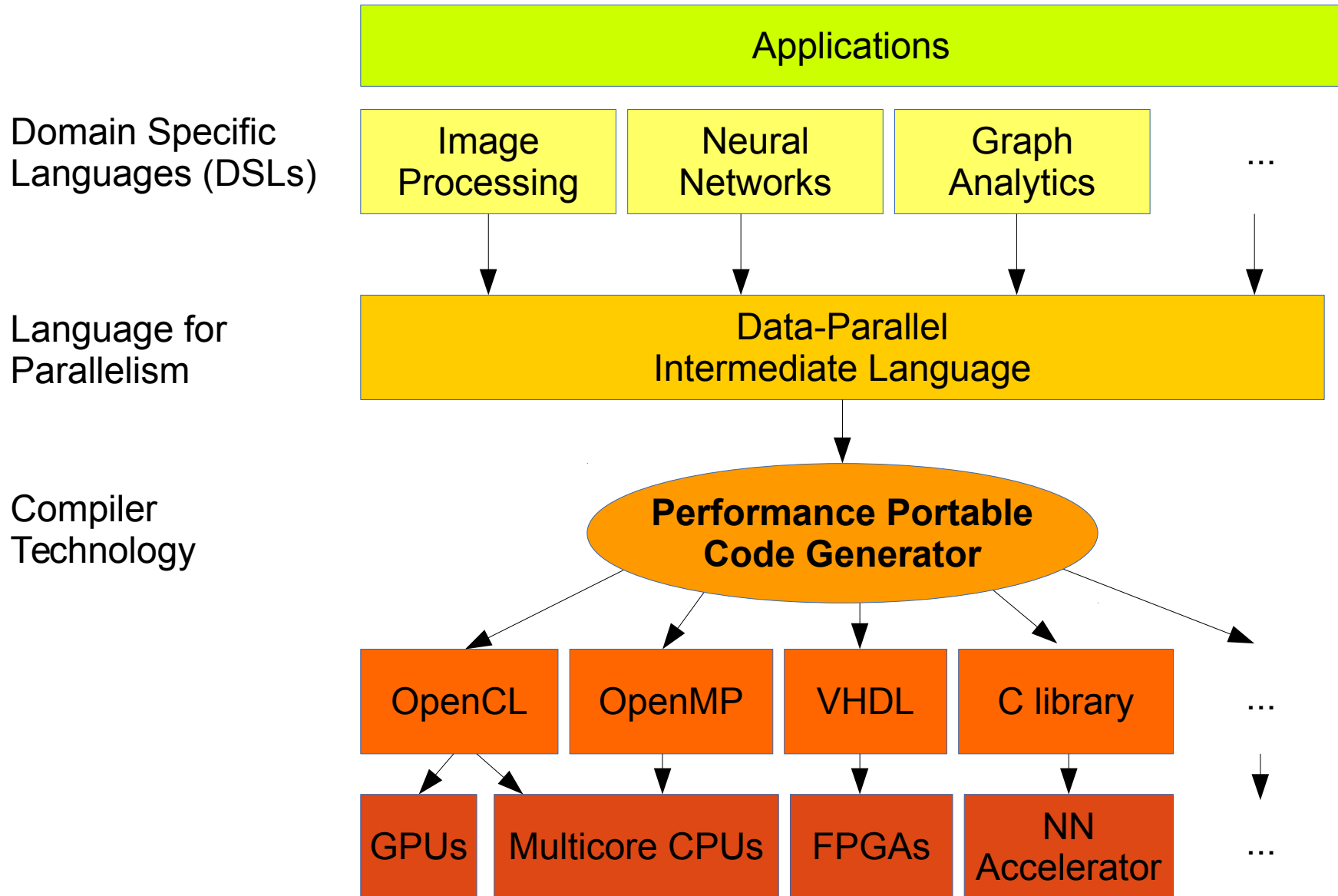
Performance is not portable



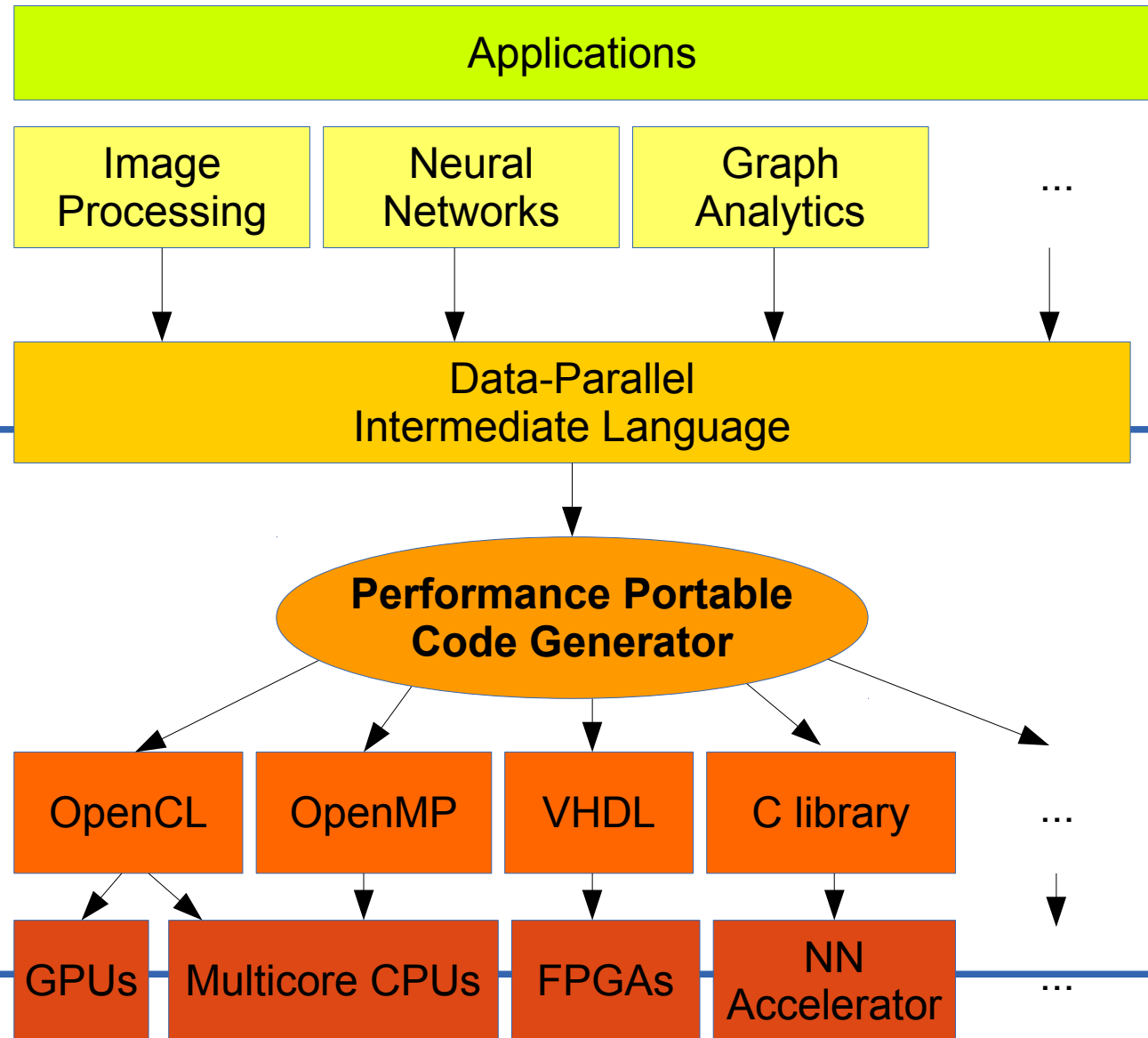
Current landscape



What we need

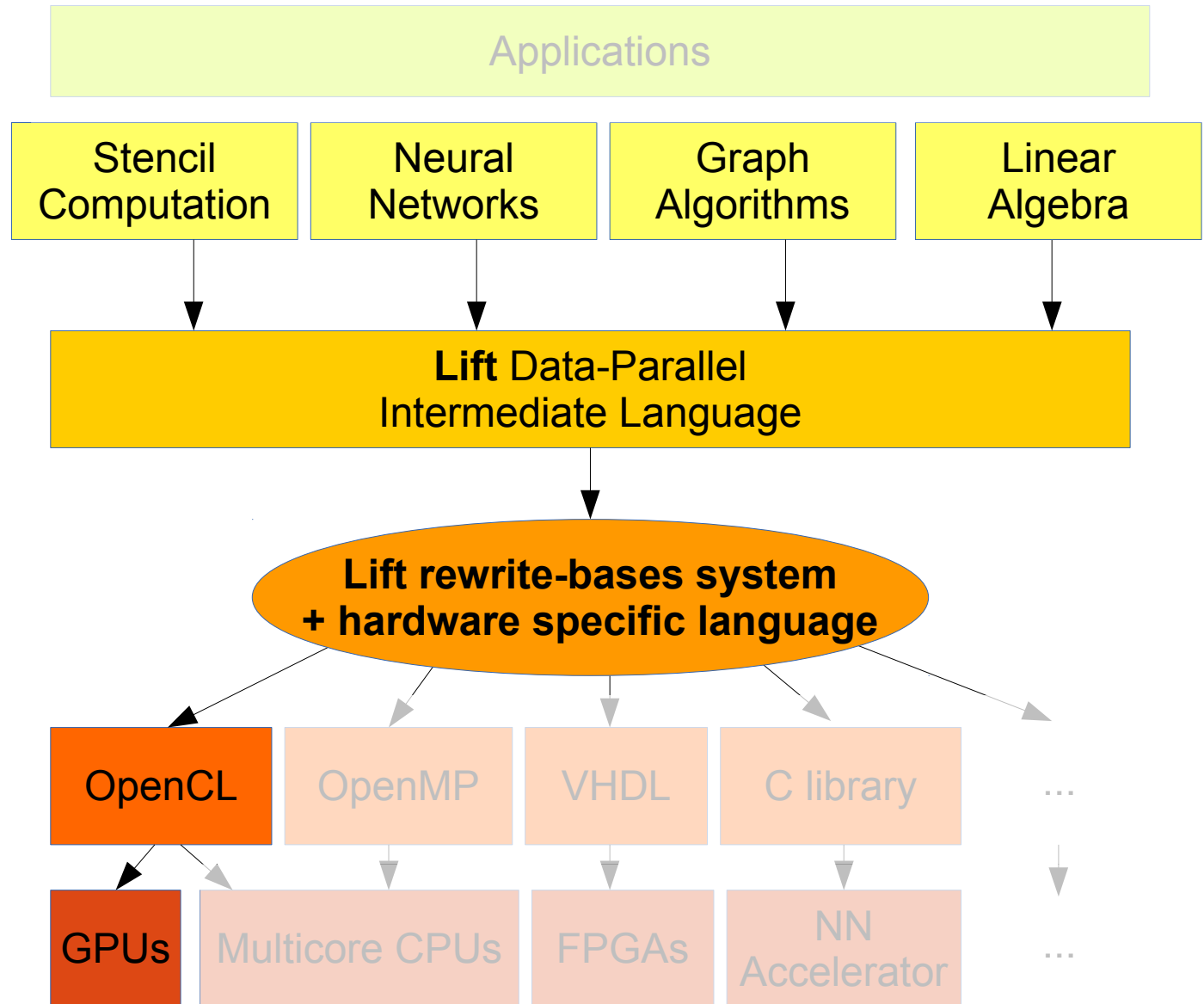


Bottom up Approach



**Lift
(bottom-up)**

Current work



Evolution of programming

Evolution of programming

Assembly

```
        la    $s0, input;
        li    $t1, 0
sum:    beq   $t1, $t0, endSun
        lw    $t7, 0(s0)
        addi $t4, $t4, $t7
        addi $s0, $s0, 4
        addi $t3, $t3, 1
        jump sum
```

Evolution of programming

Assembly

```
        la    $s0, input;
        li    $t1, 0
sum:    beq   $t1, $t0, endSun
        lw    $t7, 0(s0)
        addi $t4, $t4, $t7
        addi $s0, $s0, 4
        addi $t3, $t3, 1
        jump sum
```

Structured

```
float acc = 0;
for (int i=0; i<N; i++)
    acc += input[i];
```

Evolution of programming

Assembly

```
        la    $s0, input;
        li    $t1, 0
sum:    beq   $t1, $t0, endSun
        lw    $t7, 0(s0)
        addi $t4, $t4, $t7
        addi $s0, $s0, 4
        addi $t3, $t3, 1
        jump sum
```

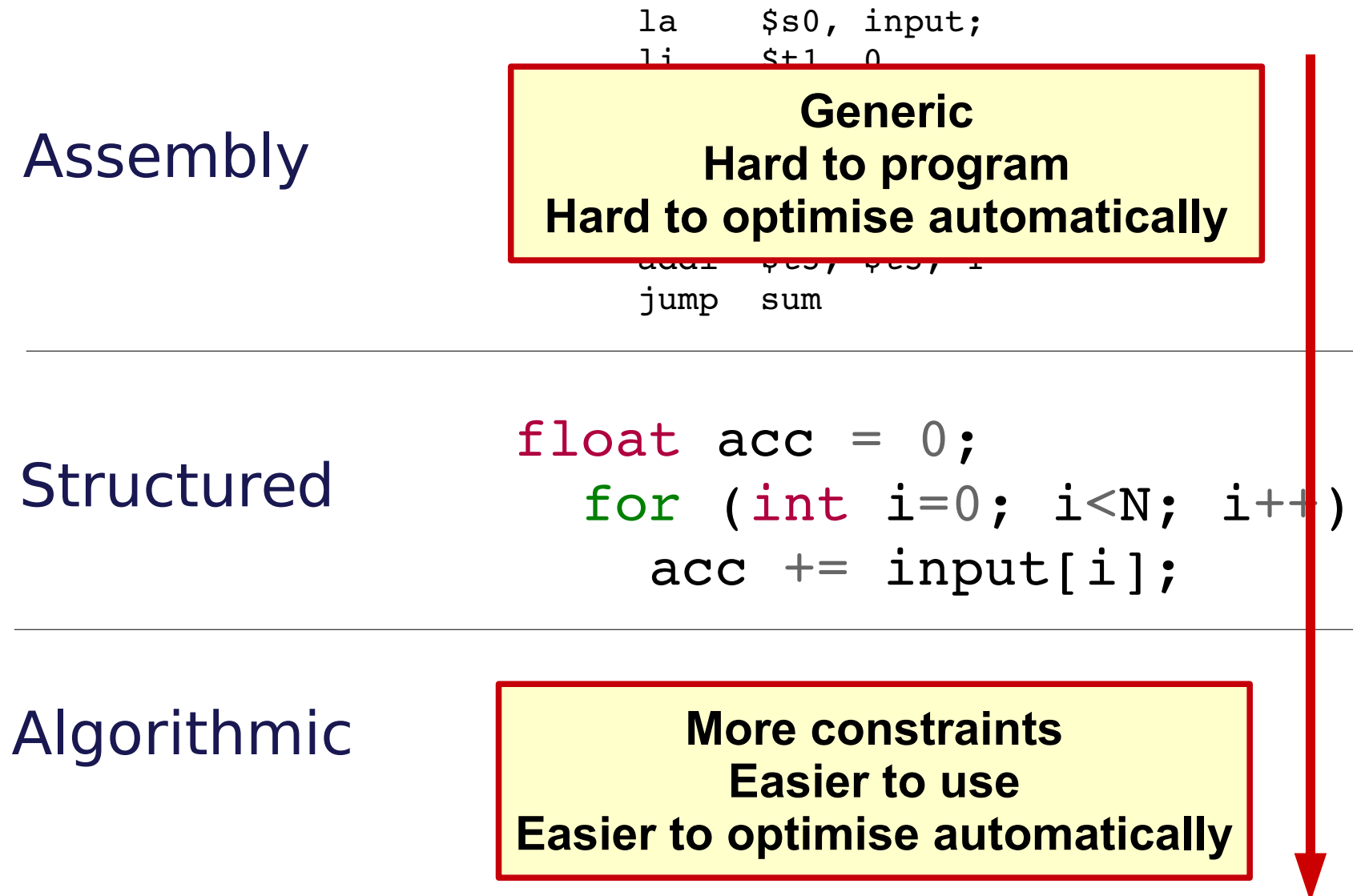
Structured

```
float acc = 0;
for (int i=0; i<N; i++)
    acc += input[i];
```

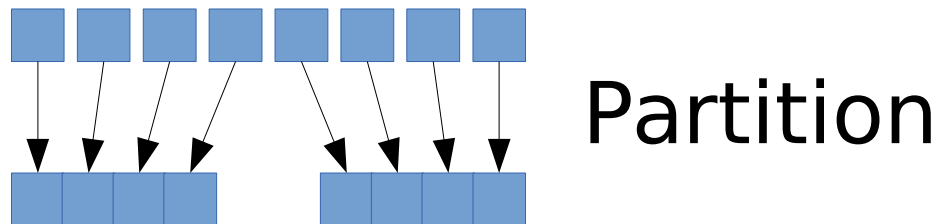
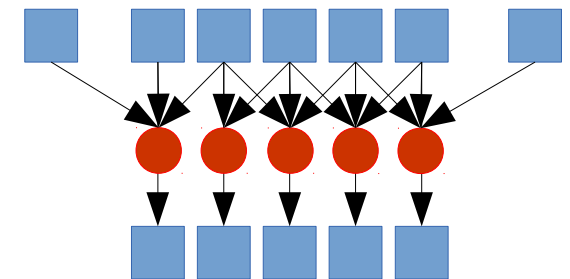
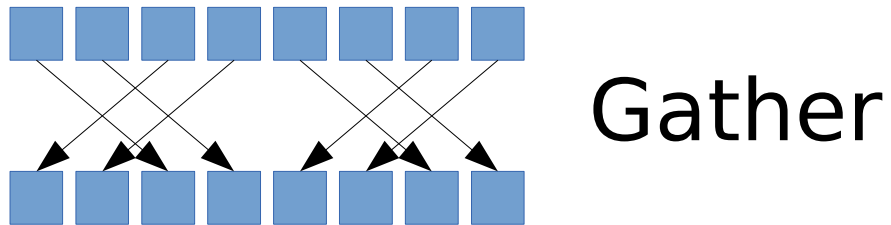
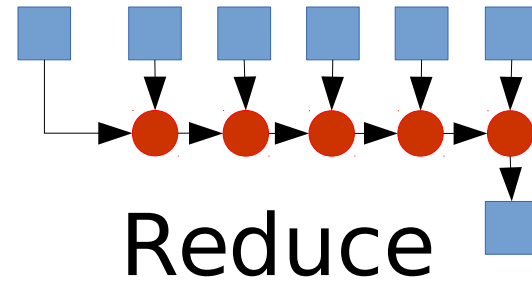
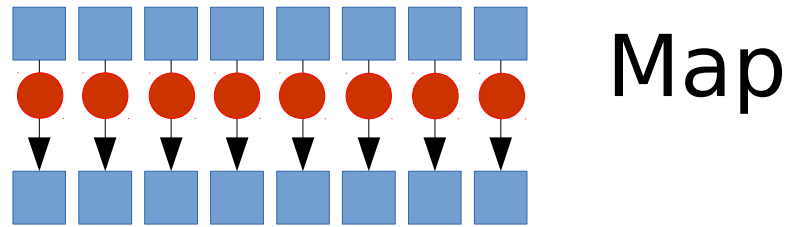
Algorithmic

```
reduce(0, +, input)
```

Evolution of programming



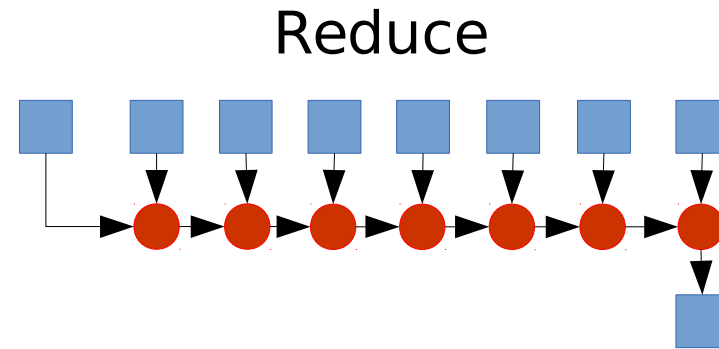
Algorithmic Patterns



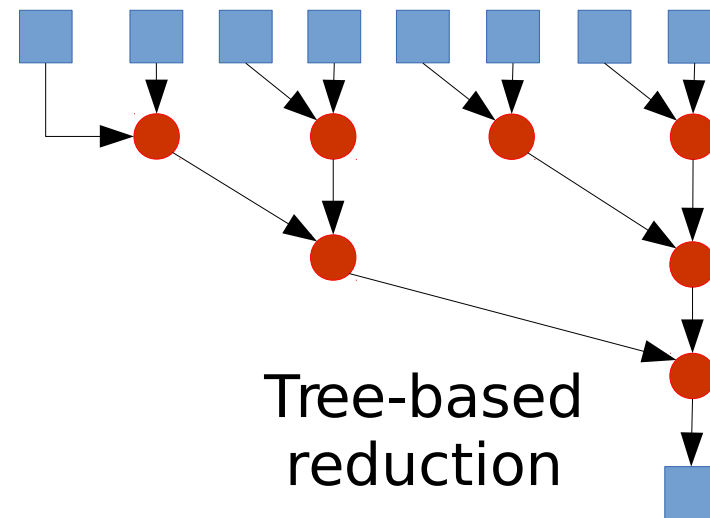
...

Separation of concern

**Programming
abstraction**



Implementation



Back to high performance code

```
kernel void sum(global float* g_in, global float* g_out,
               unsigned int n, local volatile float* l_data) {

    unsigned int tid = get_local_id(0);
    unsigned int i   = get_group_id(0) * 256 + get_local_id(0);

    l_data[tid] = 0;
    while (i < n) {
        l_data[tid] += g_in[i];
        i += 256 * get_num_groups(0);
    }
    barrier(CLK_LOCAL_MEM_FENCE);

    if (tid < 128)
        l_data[tid] += l_data[tid+128];
    barrier(CLK_LOCAL_MEM_FENCE);
    if (tid < 64)
        l_data[tid] += l_data[tid+ 64];
    barrier(CLK_LOCAL_MEM_FENCE);

    if (tid < 32) {
        l_data[tid] += l_data[tid+32]; l_data[tid] += l_data[tid+16];
        l_data[tid] += l_data[tid+ 8]; l_data[tid] += l_data[tid+ 4];
        l_data[tid] += l_data[tid+ 2]; l_data[tid] += l_data[tid+ 1];
    }
    if (tid == 0)
        g_out[get_group_id(0)] = l_data[0];
}
```



How did we get there?

- Built step by step by hand
- Human “knows” good patterns of optimisations

Back to high performance code

```
kernel void sum(global float* g_in, global float* g_out,
               unsigned int n, local volatile float* l_data) {

    unsigned int tid = get_local_id(0);
    unsigned int i   = get_group_id(0) * 256 + get_local_id(0);

    l_data[tid] = 0;
    while (i < n) {
        l_data[tid] += g_in[i];
        i += 256 * get_num_groups(0);
    }
    barrier(CLK_LOCAL_MEM_FENCE);

    if (tid < 128)
        l_data[tid] += l_data[tid+128];
    barrier(CLK_LOCAL_MEM_FENCE);
    if (tid < 64)
        l_data[tid] += l_data[tid+ 64];
    barrier(CLK_LOCAL_MEM_FENCE)

    if (tid < 32) {
        l_data[tid] += l_data[tid+32]; l_data[tid] += l_data[tid+16];
        l_data[tid] += l_data[tid+ 8]; l_data[tid] += l_data[tid+ 4];
        l_data[tid] += l_data[tid+ 2]; l_data[tid] += l_data[tid+ 1];
    }
    if (tid == 0)
        g_out[get_group_id(0)] = l_data[0];
}
```



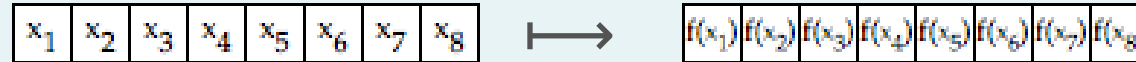
How did we get there?

How to get there automatically?

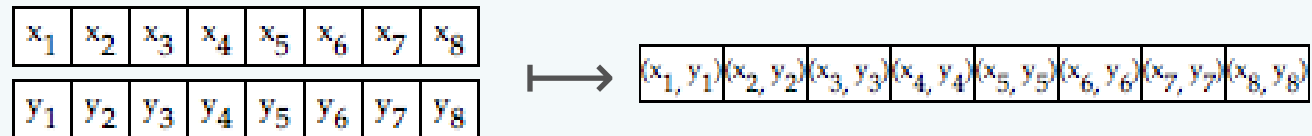
The Lift Approach

Lift Intermediate Language

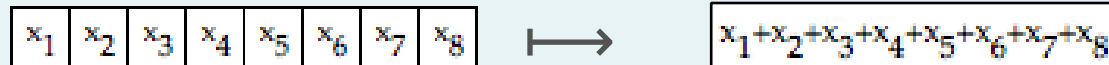
map(f) :



zip:



reduce(0,+):



split(n):



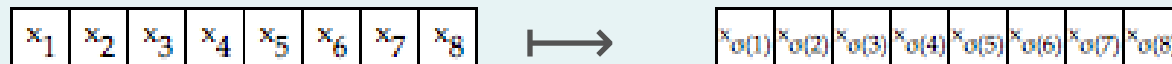
join:



iterate(f, n):



reorder(σ):



Focus on the **what** rather than **how**

Abstracting the implementation

```
kernel void sum(global float* g_in, global float* g_out,
               unsigned int n, local volatile float* l_data) {

    unsigned int tid = get_local_id(0);
    unsigned int i   = get_group_id(0) * 256 + get_local_id(0);

    l_data[tid] = 0;
    while (i < n) {
        l_data[tid] += g_in[i];
        i += 256 * get_num_groups(0);
    }
    barrier(CLK_LOCAL_MEM_FENCE);

    if (tid < 128)
        l_data[tid] += l_data[tid+128];
    barrier(CLK_LOCAL_MEM_FENCE);
    if (tid < 64)
        l_data[tid] += l_data[tid+ 64];
    barrier(CLK_LOCAL_MEM_FENCE)

    if (tid < 32) {
        l_data[tid] += l_data[tid+32]; l_data[tid] += l_data[tid+16];
        l_data[tid] += l_data[tid+ 8]; l_data[tid] += l_data[tid+ 4];
        l_data[tid] += l_data[tid+ 2]; l_data[tid] += l_data[tid+ 1];
    }
    if (tid == 0)
        g_out[get_group_id(0)] = l_data[0];
}
```

OpenCL

```
split(256) >>
reorder >>
mapWorkGroup(
    reduceSeq(
        toLocal(
            mapLocalThread(0)),
        zip(input) >>
        mapLocalThread(
            +)) >>
iterate(
    reorder >>
    split(2) >>
    mapLocalThread(
        reduceSeq(0,+))) >>
iterate(
    reorder >>
    split(2) >>
    mapWarp(
        reduceSeq(0,+))) >>
toGlobal(
    mapLocalThread(id))
```

Lift IR

Abstracting the implementation

```
kernel void sum(global float* g_in, global float* g_out,
               unsigned int n, local volatile float* l_data) {
    unsigned int tid = get_local_id(0);
    unsigned int i   = get_group_id(0) * 256 + get_local_id(0);

    l_data[tid] = 0;
    while (i < n) {
        l_data[tid] += g_in[i];
        i += 256 * get_num_groups(0);
    }
    barrier(CLK_LOCAL_MEM_FENCE);

    if (tid < 128)
        l_data[tid] += l_data[tid+128];
    barrier(CLK_LOCAL_MEM_FENCE);
    if (tid < 64)
        l_data[tid] += l_data[tid+ 64];
    barrier(CLK_LOCAL_MEM_FENCE)

    if (tid < 32) {
        l_data[tid] += l_data[tid+32]; l_data[tid] += l_data[tid+16];
        l_data[tid] += l_data[tid+ 8]; l_data[tid] += l_data[tid+ 4];
        l_data[tid] += l_data[tid+ 2]; l_data[tid] += l_data[tid+ 1];
    }
    if (tid == 0)
        g_out[get_group_id(0)] = l_data[0];
}
```

OpenCL

```
split(256) >>
reorder >>
mapWorkGroup(
    reduceSeq(
        toLocal(
            mapLocalThread(0)),
        zip(input) >>
        mapLocalThread(
            +)) >>
    iterate(
        reorder >>
        split(2) >>
        mapLocalThread(
            reduceSeq(0,+))) >>
    iterate(
        reorder >>
        split(2) >>
        mapWarp(
            reduceSeq(0,+))) >>
    toGlobal(
        mapLocalThread(id)))
```

Lift IR

Abstracting the implementation

```
kernel void sum(global float* g_in, global float* g_out,
               unsigned int n, local volatile float* l_data) {

    unsigned int tid = get_local_id(0);
    unsigned int i   = get_group_id(0) * 256 + get_local_id(0);

    l_data[tid] = 0;
    while (i < n) {
        l_data[tid] += g_in[i];
        i += 256 * get_num_groups(0);
    }
    barrier(CLK_LOCAL_MEM_FENCE);

    if (tid < 128)
        l_data[tid] += l_data[tid+128];
    barrier(CLK_LOCAL_MEM_FENCE);
    if (tid < 64)
        l_data[tid] += l_data[tid+ 64];
    barrier(CLK_LOCAL_MEM_FENCE)

    if (tid < 32) {
        l_data[tid] += l_data[tid+32]; l_data[tid] += l_data[tid+16];
        l_data[tid] += l_data[tid+ 8]; l_data[tid] += l_data[tid+ 4];
        l_data[tid] += l_data[tid+ 2]; l_data[tid] += l_data[tid+ 1];
    }
    if (tid == 0)
        g_out[get_group_id(0)] = l_data[0];
}
```

OpenCL

```
split(256) >>
reorder >>
mapWorkGroup(
    reduceSeq(
        toLocal(
            mapLocalThread(0)),
        zip(input) >>
        mapLocalThread(
            +)) >>
    iterate(
        reorder >>
        split(2) >>
        mapLocalThread(
            reduceSeq(0,+))) >>
    iterate(
        reorder >>
        split(2) >>
        mapWarp(
            reduceSeq(0,+))) >>
    toGlobal(
        mapLocalThread(id)))
```

Lift IR

Abstracting the implementation

```
kernel void sum(global float* g_in, global float* g_out,
               unsigned int n, local volatile float* l_data) {

    unsigned int tid = get_local_id(0);
    unsigned int i   = get_group_id(0) * 256 + get_local_id(0);

    l_data[tid] = 0;
    while (i < n) {
        l_data[tid] += g_in[i];
        i += 256 * get_num_groups(0);
    }
    barrier(CLK_LOCAL_MEM_FENCE);

    if (tid < 128)
        l_data[tid] += l_data[tid+128];
    barrier(CLK_LOCAL_MEM_FENCE);
    if (tid < 64)
        l_data[tid] += l_data[tid+ 64];
    barrier(CLK_LOCAL_MEM_FENCE)

    if (tid < 32) {
        l_data[tid] += l_data[tid+32]; l_data[tid] += l_data[tid+16];
        l_data[tid] += l_data[tid+ 8]; l_data[tid] += l_data[tid+ 4];
        l_data[tid] += l_data[tid+ 2]; l_data[tid] += l_data[tid+ 1];
    }
    if (tid == 0)
        g_out[get_group_id(0)] = l_data[0];
}
```

OpenCL

```
split(256) >>
reorder >>
mapWorkGroup(
    reduceSeq(
        toLocal(
            mapLocalThread(0)),
        zip(input) >>
        mapLocalThread(
            +)) >>
    iterate(
        reorder >>
        split(2) >>
        mapLocalThread(
            reduceSeq(0,+))) >>
    iterate(
        reorder >>
        split(2) >>
        mapWarp(
            reduceSeq(0,+))) >>
    toGlobal(
        mapLocalThread(id)))
```

Lift IR

Abstracting the implementation

```
kernel void sum(global float* g_in, global float* g_out,
               unsigned int n, local volatile float* l_data) {

    unsigned int tid = get_local_id(0);
    unsigned int i   = get_group_id(0) * 256 + get_local_id(0);

    l_data[tid] = 0;
    while (i < n) {
        l_data[tid] += g_in[i];
        i += 256 * get_num_groups(0);
    }
    barrier(CLK_LOCAL_MEM_FENCE);

    if (tid < 128)
        l_data[tid] += l_data[tid+128];
    barrier(CLK_LOCAL_MEM_FENCE);
    if (tid < 64)
        l_data[tid] += l_data[tid+ 64];
    barrier(CLK_LOCAL_MEM_FENCE)

    if (tid < 32) {
        l_data[tid] += l_data[tid+32]; l_data[tid] += l_data[tid+16];
        l_data[tid] += l_data[tid+ 8]; l_data[tid] += l_data[tid+ 4];
        l_data[tid] += l_data[tid+ 2]; l_data[tid] += l_data[tid+ 1];
    }
    if (tid == 0)
        g_out[get_group_id(0)] = l_data[0];
}
```

OpenCL

```
split(256) >>
reorder >>
mapWorkGroup(
    reduceSeq(
        toLocal(
            mapLocalThread(0)),
        zip(input) >>
        mapLocalThread(
            +)) >>
iterate(
    reorder >>
    split(2) >>
    mapLocalThread(
        reduceSeq(0,+))) >>
iterate(
    reorder >>
    split(2) >>
    mapWarp(
        reduceSeq(0,+))) >>
toGlobal(
    mapLocalThread(id))
```

Lift IR

Abstracting the implementation

```
kernel void sum(global float* g_in, global float* g_out,
               unsigned int n, local volatile float* l_data) {

    unsigned int tid = get_local_id(0);
    unsigned int i   = get_group_id(0) * 256 + get_local_id(0);

    l_data[tid] = 0;
    while (i < n) {
        l_data[tid] += g_in[i];
        i += 256 * get_num_groups(0);
    }
    barrier(CLK_LOCAL_MEM_FENCE);

    if (tid < 128)
        l_data[tid] += l_data[tid+128];
    barrier(CLK_LOCAL_MEM_FENCE);
    if (tid < 64)
        l_data[tid] += l_data[tid+ 64];
    barrier(CLK_LOCAL_MEM_FENCE)

    if (tid < 32) {
        l_data[tid] += l_data[tid+32]; l_data[tid] += l_data[tid+16];
        l_data[tid] += l_data[tid+ 8]; l_data[tid] += l_data[tid+ 4];
        l_data[tid] += l_data[tid+ 2]; l_data[tid] += l_data[tid+ 1];
    }
    if (tid == 0)
        g_out[get_group_id(0)] = l_data[0];
}
```

OpenCL

```
split(256) >>
reorder >>
mapWorkGroup(
    reduceSeq(
        toLocal(
            mapLocalThread(0)),
        zip(input) >>
        mapLocalThread(
            +)) >>
iterate(
    reorder >>
    split(2) >>
    mapLocalThread(
        reduceSeq(0,+))) >>
iterate(
    reorder >>
    split(2) >>
    mapWarp(
        reduceSeq(0,+))) >>
toGlobal(
    mapLocalThread(id))
```

Lift IR

Observations

- ▶ Can describe implementation as composition of primitives
 - map, reduce, zip, ...
 - same ones we use at the high level!
- ▶ Can express complex optimisations
 - e.g. shared memory use

How to generate optimised version?

Programming abstraction

`reduce(0,+,input)`



Implementation

```
split(256) >>
reorder >>
mapWorkGroup(
  reduceSeq(
    toLocal(
      mapLocalThread(0)),
    zip(input) >>
    mapLocalThread(
      +)) >>
  iterate(
    reorder >>
    split(2) >>
    mapLocalThread(
      reduceSeq(0,+)) >>
  iterate(
    reorder >>
    split(2) >>
    mapWarp(
      reduceSeq(0,+)) >>
  toGlobal(
    mapLocalThread(id)))
```

Algorithmic Rewrite Rules

- Provably correct rewrite rules
- Express algorithmic implementation choices

Split-join rule:

`map(f) → split(n) >> map(map(f)) >> join`

Map fusion rule:

`map(f) >> map(g) → map(f >> g)`

Reduce rules:

`reduce(z, f) → reducePart(z, f) >> reduce(f)`

`reducePart(z, f) → reorder >> reducePart(z, f)`

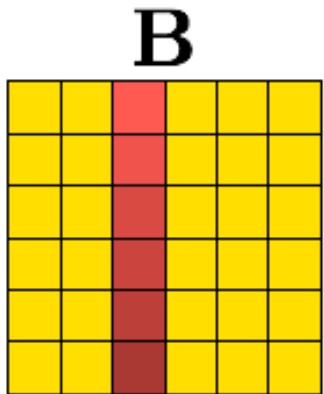
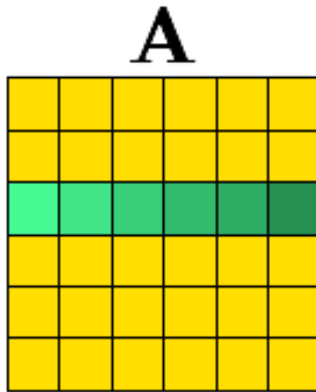
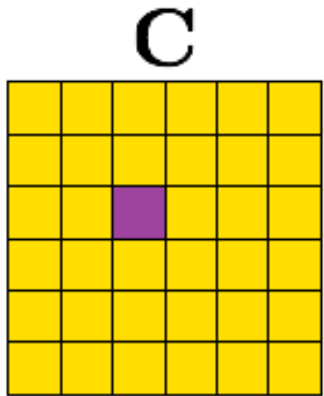
`reducePart(z, f) → split(n) >> map(reducePart(z, f)) >> join`

`reducePart(z, f) → iterate(reducePart(z, f))`

`reducePart(z, f) → reduceSeq(z, f)`

Matrix-multiplication example

Matrix Multiplication in Lift



```
A >> map(λ rowOfA ↦  
B >> map(λ colOfB ↦  
  zip(rowOfA, colOfB) >>  
  map(mult) >> reduce(0.0f, add)  
)  
)
```

```
A >> map( $\lambda$  rowOfA  $\mapsto$ 
  B >> map( $\lambda$  colOfB  $\mapsto$ 
    zip(rowOfA, colOfB) >>
    map(mult) >> reduce(0.0f, add)
  )
)
```

```
A >> map( $\lambda$  rowOfA  $\mapsto$ 
  B >> map( $\lambda$  colOfB  $\mapsto$ 
    zip(rowOfA, colOfB) >>
    map(mult) >> reduce(0.0f, add)
  )
)
```



```
1 for (int i = 0; i < M; i++) {
2   for (int j = 0; j < N; j++) {
3     for (int k = 0; k < K; k++) {
4       temp[k + K*N i + K*j] =
5         mult(A[k + K*i], B[k + K*j]);
6     }
7   for (int k = 0; k < K; k++) {
8     C[j + N*i] +=
9       temp[k + K*N*i + K*j];
10  }
11 }
12 }
```

```

A >> map (λ rowOfA ↦
  B >> map(λ colOfB ↦
    zip(rowOfA, colOfB) >>
      map(mult) >> reduce(0.0f, add)
  )
)

```



```

1 for (int i = 0; i < M; i++) {
2   for (int j = 0; j < N; j++) {
3     for (int k = 0; k < K; k++) {
4       temp[k + K*N i + K*j] =
5         mult(A[k + K*i], B[k + K*j]);
6     }
7   for (int k = 0; k < K; k++) {
8     C[j + N*i] +=
9       temp[k + K*N*i + K*j];
10  }
11 }
12 }

```

Map(f) ⇒ Join() ∘ Map(Map(f)) ∘ Split(k)

```

A >> map (λ rowOfA ↦
  B >> map(λ colOfB ↦
    zip(rowOfA, colOfB) >>
    map(mult) >> reduce(0.0f, add)
  )
)

```

```

1 for (int i = 0; i < M; i++) {
2   for (int j = 0; j < N; j++) {
3     for (int k = 0; k < K; k++) {
4       temp[k + K*N i + K*j] =
5         mult(A[k + K*i], B[k + K*j]);
6     }
7   for (int k = 0; k < K; k++) {
8     C[j + N*i] +=
9       temp[k + K*N*i + K*j];
10  }
11 }
12 }

```

Map(f) ⇒ Join() ∘ Map(Map(f)) ∘ Split(k)

```

A >> split(m) >> map (λ rowsOfA ↦
  rowsOfA >> map (λ rowOfA ↦
    B >> map(λ colOfB ↦
      zip(rowOfA, colOfB) >>
      map(mult) >> reduce(0.0f, add)
    )
  )
) >> join

```

```

A >> map (λ rowOfA ↦
  B >> map(λ colOfB ↦
    zip(rowOfA, colOfB) >>
    map(mult) >> reduce(0.0f, add)
  )
)

```

```

1 for (int i = 0; i < M; i++) {
2   for (int j = 0; j < N; j++) {
3     for (int k = 0; k < K; k++) {
4       temp[k + K*N i + K*j] =
5         mult(A[k + K*i], B[k + K*j]);
6     }
7   }
8   for (int k = 0; k < K; k++) {
9     C[j + N*i] +=
10      temp[k + K*N*i + K*j];
11   }
12 }

```

Map(f) ⇒ Join() ∘ Map(Map(f)) ∘ Split(k)

```

A >> split(m) >> map (λ rowsOfA ↦
  rowsOfA >> map (λ rowOfA ↦
    B >> map(λ colOfB ↦
      zip(rowOfA, colOfB) >>
      map(mult) >> reduce(0.0f, add)
    )
  )
) >> join

```

```

1 for (int i = 0; i < M/m; i++) {
2   for (int l = 0; l < m; l++) {
3     for (int j = 0; j < N; j++) {
4       for (int k = 0; k < K; k++) {
5         temp[k + 2*K*N*i + K*N*l + K*j] =
6           mult(A[k + K*l + 2*K*i], B[k + K*j]);
7       }
8     }
9     for (int k = 0; k < K; k++) {
10      C[j + N*l + 2*N*i] +=
11       temp[k + 2*K*N*i + K*N*l + K*j];
12    }
13  }
14 }

```

Map interchanged

```
A >> split(m) >> map( $\lambda$  rowsOfA  $\mapsto$ 
  B >> map( $\lambda$  colOfB  $\mapsto$ 
    rowsOfA >> map( $\lambda$  rowOfA  $\mapsto$ 
      zip(rowOfA, colOfB) >>
        map(mult) >> reduce(0.0f, add)
    )
  ) >> transpose
) >> join
```



```
1 for (int i = 0; i < M/2; i++) {
2   for (int j = 0; j < N; j++) {
3     for (int l = 0; l < 2; l++) {
4       for (int k = 0; k < K; k++) {
5         temp[k + 2*K*N*i + K*N*l + K*j] =
6           mult(A[k + K*l + 2*K*i], B[k + K*j]);
7       }
8     }
9     for (int k = 0; k < K; k++) {
10      C[j + N*l + 2*N*i] +=
11        temp[k + 2*K*N*i + K*N*l + K*j];
12    }
13  }
14 }
```

Split-join rule

```
A >> split(m) >> map(λ rowsOfA ↦
  B >> split(n) >> map(λ colsOfB ↦
    colsOfB >> map(λ colOfB ↦
      rowsOfA >> map(λ rowOfA ↦
        zip(rowOfA, colOfB) >>
          map(mult) >> reduce(0.0 f, add)
      )
    )
  ) >> join >> transpose
) >> join
```



```
1 for (int i = 0; i < M/2; i++) {
2   for (int j = 0; j < N/2; j++) {
3     for (int m = 0; m < 2; m++) {
4       for (int l = 0; l < 2; l++) {
5         for (int k = 0; k < K; k++) {
6           temp[k + 4*K*N*i + 2*K*N*l + 2*K*j
7             + K*m] =
8             mult(A[k + K*l + 2*K*i], B[k + K*
9               m + 2*K*j]);
10          }
11          C[m + 2*j + 2*N*l + 4*N*i] +=
12            temp[k + 4*K*N*i + 2*K*N*l + 2*
13              K*j + K*m];
14        }
15      }
16    }
```


Map interchanged

```
A >> split(m) >> map(λ rowsOfA ↦
  B >> split(n) >> map(λ colsOfB ↦
    rowsOfA >> map(λ rowOfA ↦
      colsOfB >> map(λ colOfB ↦
        zip(rowOfA, colOfB) >>
          map(mult) >> reduce(0.0f, add)
      )
    ) >> transpose
  ) >> join >> transpose
) >> join
```



```
1 for (int i = 0; i < M/2; i++) {
2   for (int j = 0; j < N/2; j++) {
3     for (int l = 0; l < 2; l++) {
4       for (int m = 0; m < 2; m++) {
5         for (int k = 0; k < K; k++) {
6           temp[k + 4*K*N*i + 2*K*N*l + 2*K*j
7             + K*m] =
8             mult(A[k + K*l + 2*K*i], B[k + K*
9               m + 2*K*j]);
10          C[m + 2*j + 2*N*l + 4*N*i] +=
11            temp[k + 4*K*N*i + 2*K*N*l + 2*
12              K*j + K*m];
13        }
14      }
15    }
16  }
```

A few rewrite steps later...

Tiled version

```
λ (A, B) ↦
  A >> split(m) >> map(λ nRowsOfA ↦
    B >> split(n) >> map(λ mColsOfB ↦
      zip( transpose(nRowsOfA) >> split(k),
          transpose(mColsOfB) >> split(k) ) >>
      reduceSeq( init = make2DArray(n,m, 0.0 f),
        λ (accTile, (tileOfA, tileOfB)) ↦
          zip(accTile, transpose(tileOfA)) >>
          map(λ (accRow, rowOfTileOfA) ↦
            zip(accRow, transpose(tileOfB)) >>
            map(λ (acc, colOfTileOfB) ↦
              zip(rowOfTileOfA, colOfTileOfB) >>
              map(mult) >> reduce(acc, add)
            ) >> join
          )
        ) >> transpose() >>
      map(transpose) >> transpose
    ) >> join >> transpose
  ) >> join
```

```
1 for (int i = 0; i < M/2; i++) {
2   for (int j = 0; j < N/2; j++) {
3     for (int k = 0; k < K/4; k++) {
4       for (int l = 0; l < 2; l++) {
5         for (int m = 0; m < 2; m++) {
6           for (int n = 0; n < 4; n++) {
7             temp[n + 4*m + 8*N*i + 16*j + 8*l] =
8               mult(
9                 A[n + 2*K*i + 4*k + K*l],
10                B[n + 2*K*j + 4*k + K*m]
11              );
12           }
13         for (int n = 0; n < 4; n++) {
14           C[m + 2*N*i + 2*j + N*l] +=
15             temp[n + 4*m + 8*N*i + 16*j + 8*l];
16         }
17       }
18     }
19   }
20 }
21 }
```

Vectorisation

```
λ (A, B) ↦
  A >> split(m) >> map(λ nRowsOfA ↦
    B >> split(n) >> map(λ mColsOfB ↦
      zip( transpose(nRowsOfA) >> split(k),
           transpose(mColsOfB) >> split(k) ) >>
      reduceSeq( init = make2DArray(n,m, 0.0f),
                 λ (accTile, (tileOfA, tileOfB)) ↦
                   zip(accTile, transpose(tileOfA)) >>
                   map(λ (accRow, rowOfTileOfA) ↦
                     zip(accRow, transpose(tileOfB)) >>
                     map(λ (acc, colOfTileOfB) ↦
                       zip(rowOfTileOfA >> asVector(k),
                            colOfTileOfB >> asVector(k)) >>
                       map(mult4) >> asScalar >>
                       reduce(acc, add)
                     ) >> join
                   ) >> transpose() >>
                   map(transpose) >> transpose
                 ) >> join >> transpose
            ) >> join
```



```
1 for (int i = 0; i < M/2; i++) {
2   for (int j = 0; j < N/2; j++) {
3     for (int k = 0; k < K/4; k++) {
4       for (int l = 0; l < 2; l++) {
5         for (int m = 0; m < 2; m++) {
6           float4 t = mult4(
7             vload4(A, K*i/2 + k + K*l/4),
8             vload4(B, K*j/2 + k + K*m/4)
9           );
10          vstore4(t, temp, m + 2*N*i + 4*j + 2*l);
11          for (int n = 0; n < 4; n++) {
12            C[m + 2*N*i + 2*j + N*l] +=
13              temp[n + 4*m + 8*N*i + 16*j + 8*l];
14          }
15        }
16      }
17    }
18  }
19 }
```

Mapping parallelism to global threads

```
 $\lambda$  (A, B)  $\mapsto$   
A >> split(m) >> mapGlb0( $\lambda$  nRowsOfA  $\mapsto$   
B >> split(n) >> mapGlb1( $\lambda$  mColsOfB  $\mapsto$   
zip( transpose(nRowsOfA) >> split(k),  
      transpose(mColsOfB) >> split(k) ) >>  
reduceSeq( init = make2DArray(n,m, 0.0f),  
            $\lambda$  (accTile, (tileOfA, tileOfB))  $\mapsto$   
           zip(accTile, transpose(tileOfA)) >>  
           mapSeq( $\lambda$  (accRow, rowOfTileOfA)  $\mapsto$   
           zip(accRow, transpose(tileOfB)) >>  
           mapSeq( $\lambda$  (acc, colOfTileOfB)  $\mapsto$   
           zip(rowOfTileOfA >> asVector(k),  
              colOfTileOfB >> asVector(k)) >>  
           mapSeq(mult4) >> asScalar >>  
           reduceSeq(acc, add)  
           ) >> join  
           )  
           ) >> transpose() >>  
           map(transpose) >> transpose  
           ) >> join >> transpose  
           ) >> join
```



```
1 int i = get_global_id(0);  
2 int j = get_global_id(1);  
3 for (int k = 0; k < K/4; k++) {  
4     for (int l = 0; l < 2; l++) {  
5         for (int m = 0; m < 2; m++) {  
6             float4 t = mult4(  
7                 vload4(A, K*i/2 + k + K*l/4),  
8                 vload4(B, K*j/2 + k + K*m/4)  
9             );  
10            vstore4(t, temp, m + 2*N*i + 4*j + 2*l);  
11            for (int n = 0; n < 4; n++) {  
12                C[m + 2*N*i + 2*j + N*l] +=  
13                    temp[n + 4*m + 8*N*i + 16*j + 8*l];  
14            }  
15        }  
16    }  
17 }
```

Accumulating in private memory

```

λ (A, B) ↦
  A >> split(m) >> mapGlb0(λ nRowsOfA ↦
    B >> split(n) >> mapGlb1(λ mColsOfB ↦
      zip( transpose(nRowsOfA) >> split(k),
           transpose(mColsOfB) >> split(k) ) >>
      reduceSeq( init = make2DArray(n,m, 0.0 f) >>
                 toPrivate( mapSeq( mapSeq( id ) ) ),
        λ (accTile, (tileOfA, tileOfB)) ↦
          zip(accTile, transpose(tileOfA)) >>
          mapSeq(λ (accRow, rowOfTileOfA) ↦
            zip(accRow, transpose(tileOfB)) >>
            mapSeq(λ (acc, colOfTileOfB) ↦
              zip(rowOfTileOfA >> asVector(k),
                  colOfTileOfB >> asVector(k)) >>
              mapSeq(mult4) >> asScalar >>
              reduceSeq(acc, add)
            ) >> join
          )
        ) >> toGlobal( mapSeq( mapSeq( mapSeq( id ) ) ) )
    >> transpose() >>
    map(transpose) >> transpose
  ) >> join >> transpose
) >> join
  
```



```

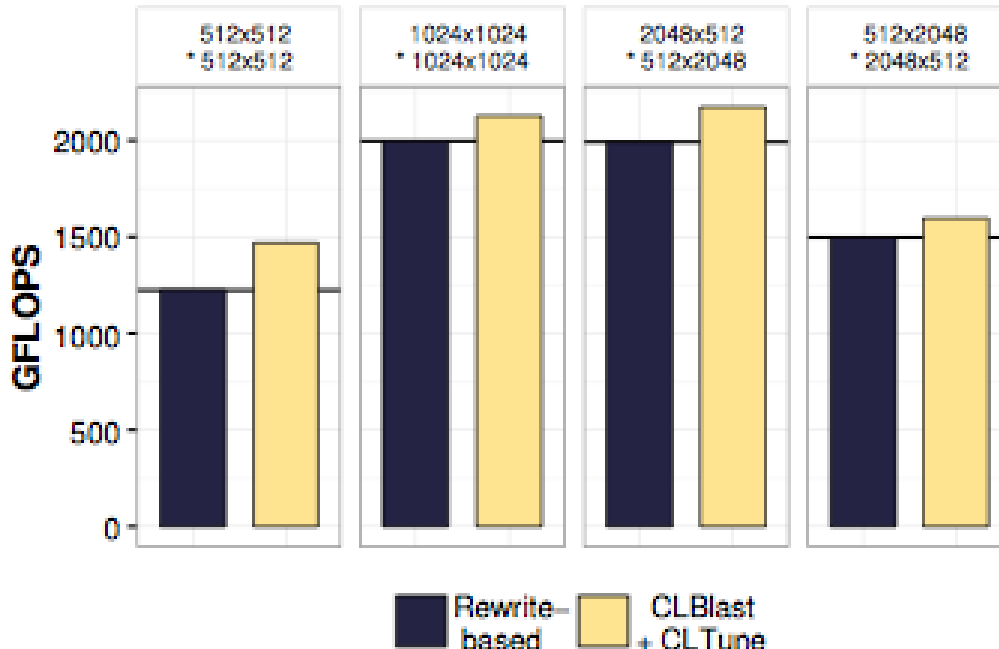
1 int i = get_global_id(0);
2 int j = get_global_id(1);
3
4 float4 temp_0; float4 temp_1;
5 float4 temp_2; float4 temp_3;
6 float acc_0; float acc_1;
7 float acc_2; float acc_3;
8
9 for (int k = 0; k < K/4; k++) {
10
11   temp_0 = mult4(vload4(k + K*i/2,A),
12                 vload4(k + K*j/2,B));
13   acc_0 += temp_0.s0 + temp_0.s1 +
14           temp_0.s2 + temp_0.s3;
15
16   temp_1 = mult4(vload4(k + K*i/2,A),
17                 vload4(k + K + 2*K*j/4,B));
18   acc_1 += temp_1.s0 + temp_1.s1 +
19           temp_1.s2 + temp_1.s3;
20
21   temp_2 = mult4(vload4(k + K + 2*K*i/4,A),
22                 vload4(k + K*j/2,B));
23   acc_2 += temp_2.s0 + temp_2.s1 +
24           temp_2.s2 + temp_2.s3;
25
26   temp_3 = mult4(vload4(k + K + 2*K*i/4, A),
27                 vload4(k + K + 2*K*j/4, B));
28   acc_3 += temp_3.s0 + temp_3.s1 +
29           temp_3.s2 + temp_3.s3;
30 }
31 C[2*N*i + 2*j] = id(acc_0);
32 C[1 + 2*N*i + 2*j] = id(acc_1);
33 C[N + 2*N*i + 2*j] = id(acc_2);
34 C[1 + N + 2*N*i + 2*j] = id(acc_3);
  
```

Performance Portability Achieved

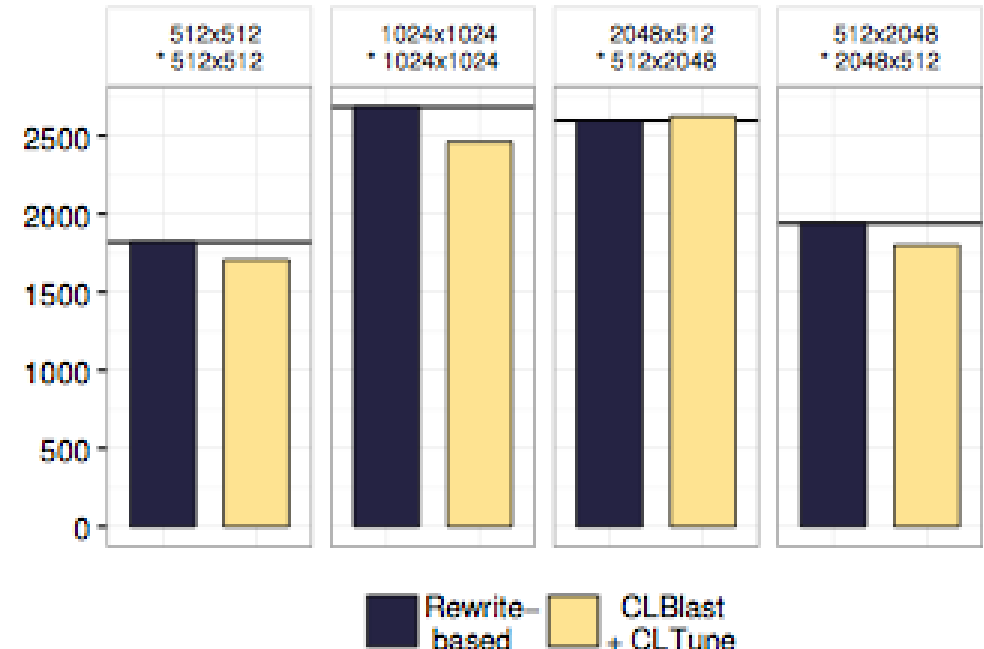
Compiler input:

```
A >> map( $\lambda$  rowOfA  $\mapsto$   
B >> map( $\lambda$  colOfB  $\mapsto$   
zip(rowOfA, colOfB) >>  
map(mult) >> reduce(0.0f, add)  
)  
)
```

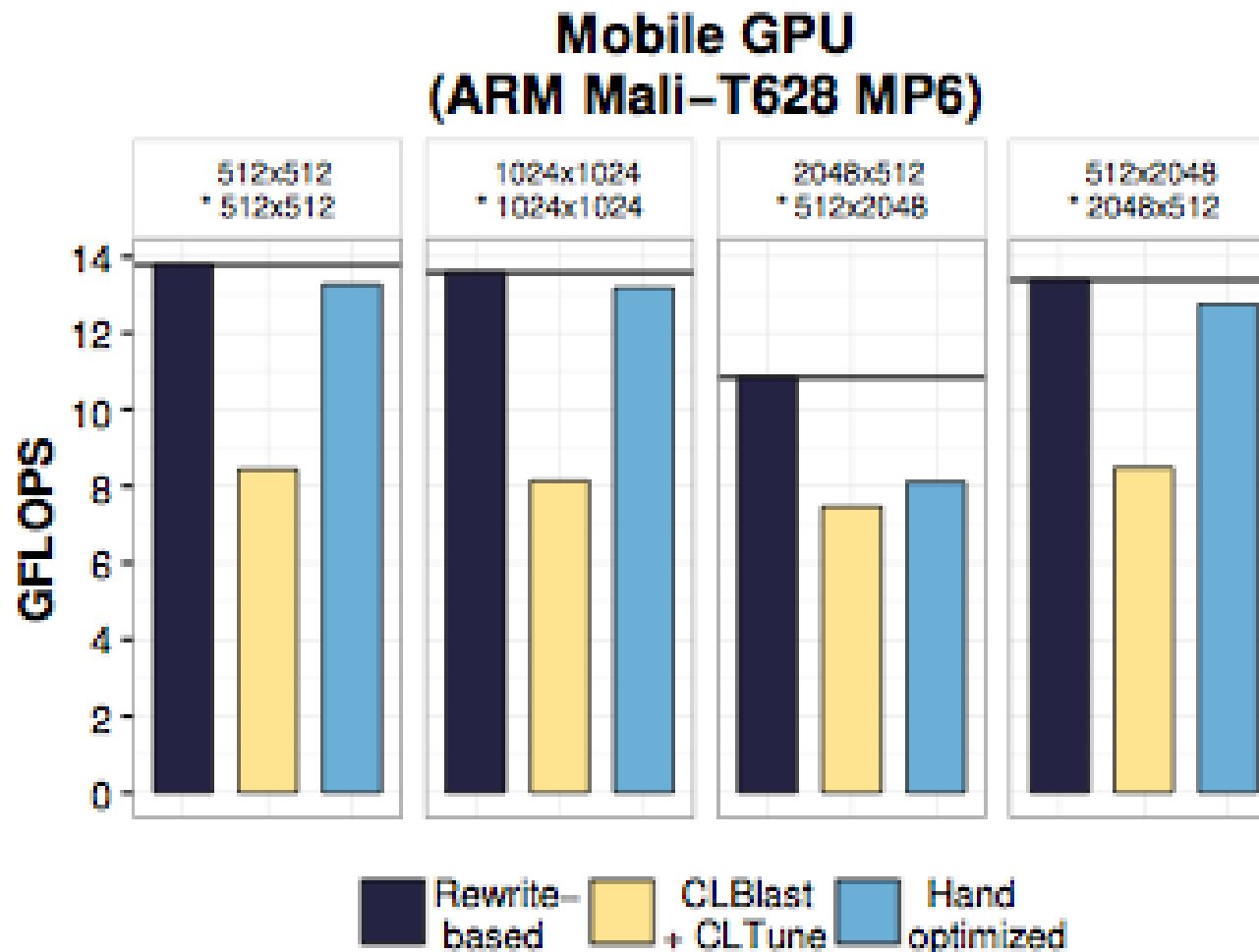
**Desktop GPU
(Nvidia GeForce GTX Titan Black)**



**Desktop GPU
(AMD Radeon HD 7970)**

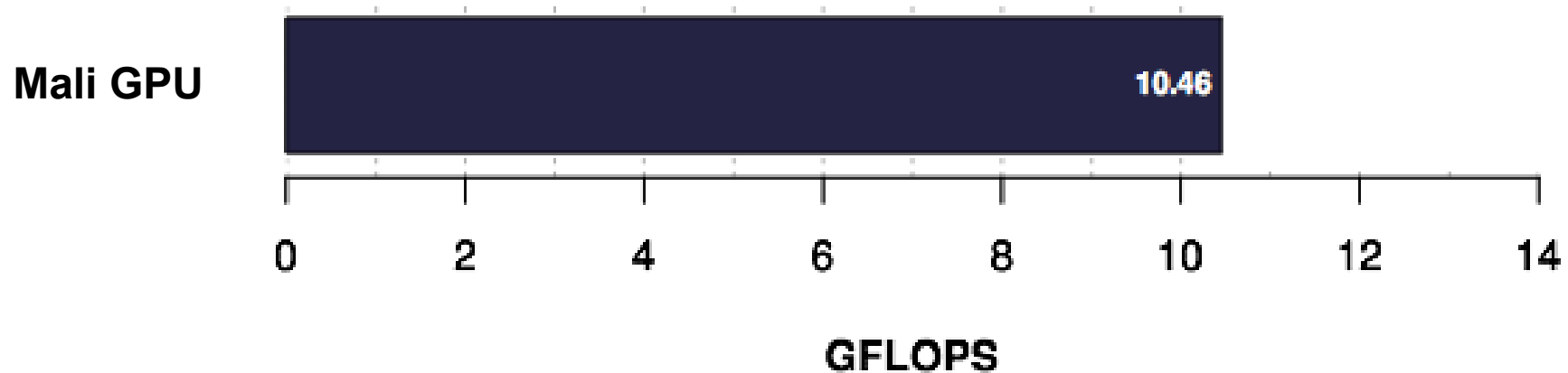


Also works on mobile GPUs



Easily Extensible

- ▶ New rules can be added
- ▶ E.g. dot built-in

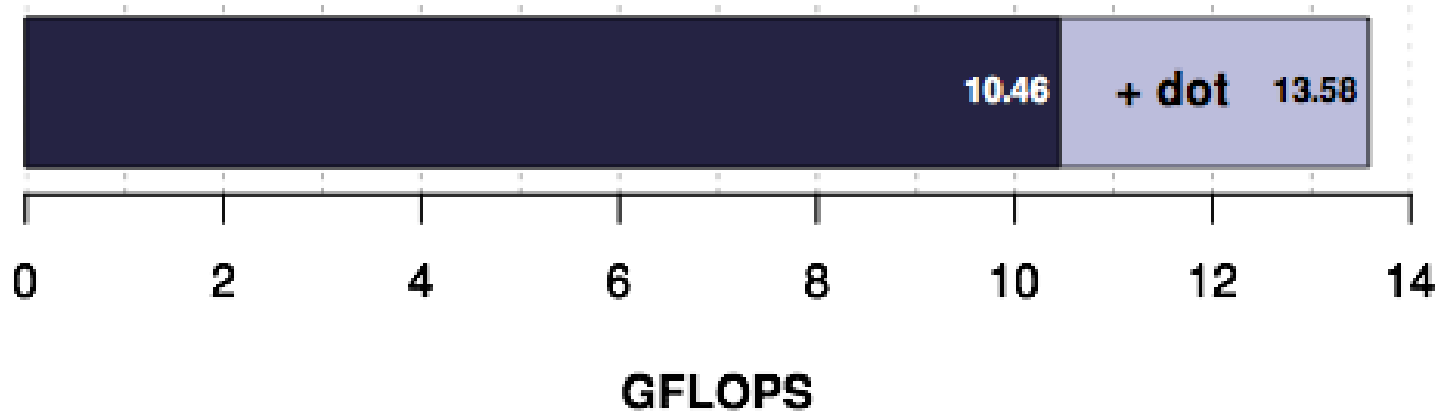


Easily Extensible

- New rules can be added
- E.g. dot built-in

```
zip(x, y) >> mapSeq(mult4) >> reduceSeq(z, add4) → dot(x, y)
```

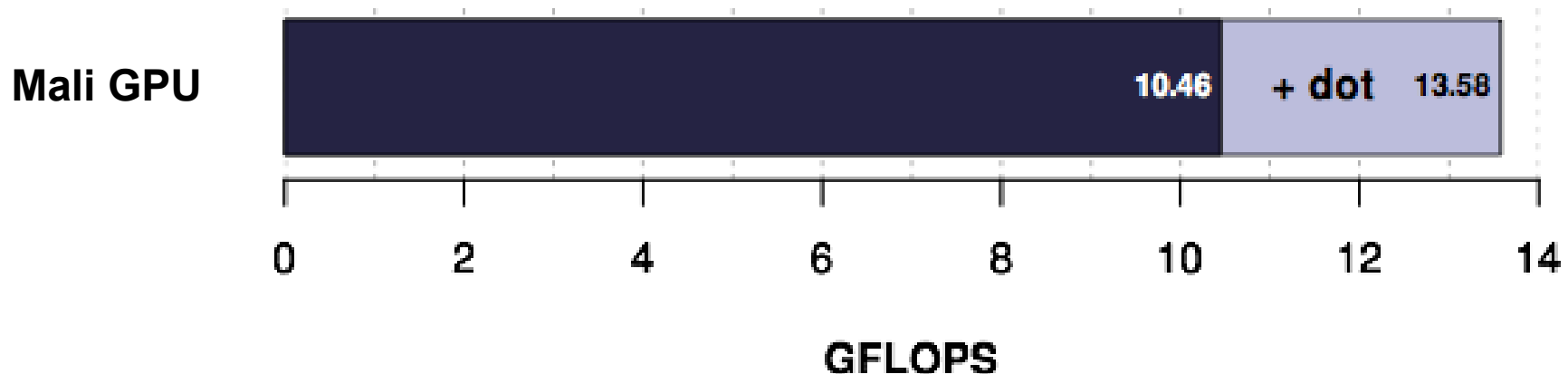
Mali GPU



Easily Extensible

- ▶ New rules can be added
- ▶ E.g. dot built-in

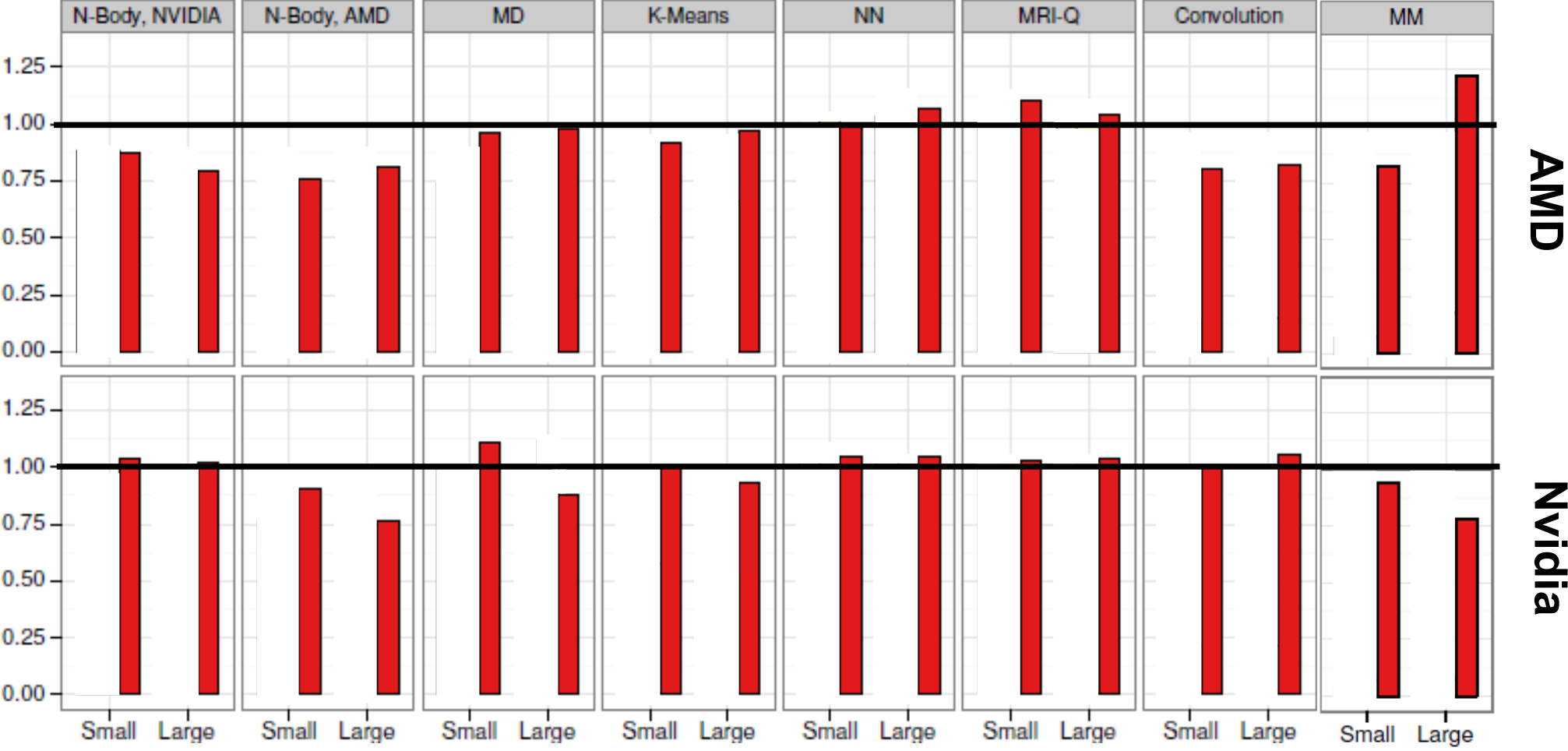
```
zip(x, y) >> mapSeq(mult4) >> reduceSeq(z, add4) → dot(x, y)
```



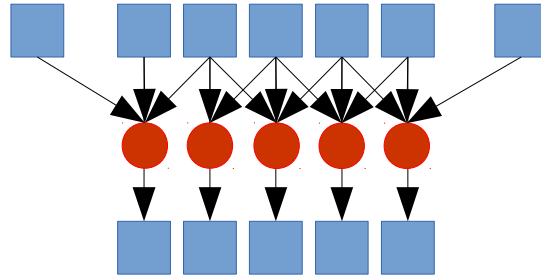
No need to modify original application!

=> Performance portable

Works for other programs too



Stencil Computation

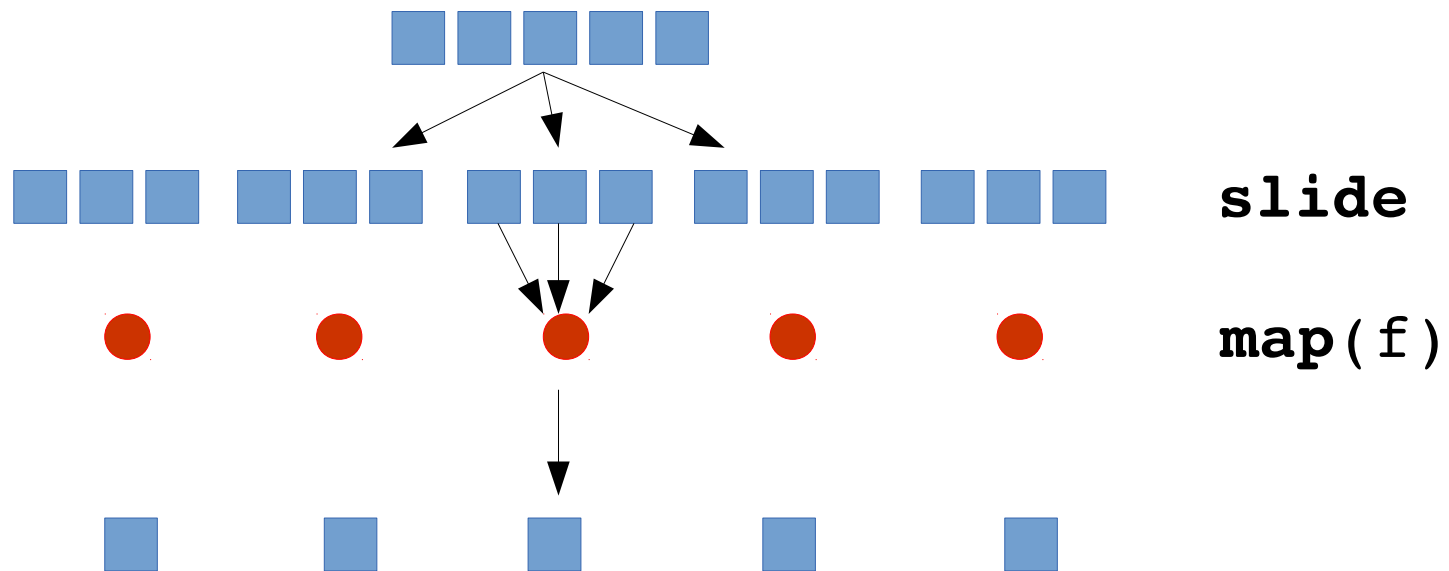


- ▶ important application
 - image processing
 - physic simulations
 - neural networks
 - pde solvers
- ▶ large opportunity for parallelism
 - GPUs perfect fit

Stencil Computation in Lift

- ▶ use a “slide” primitive

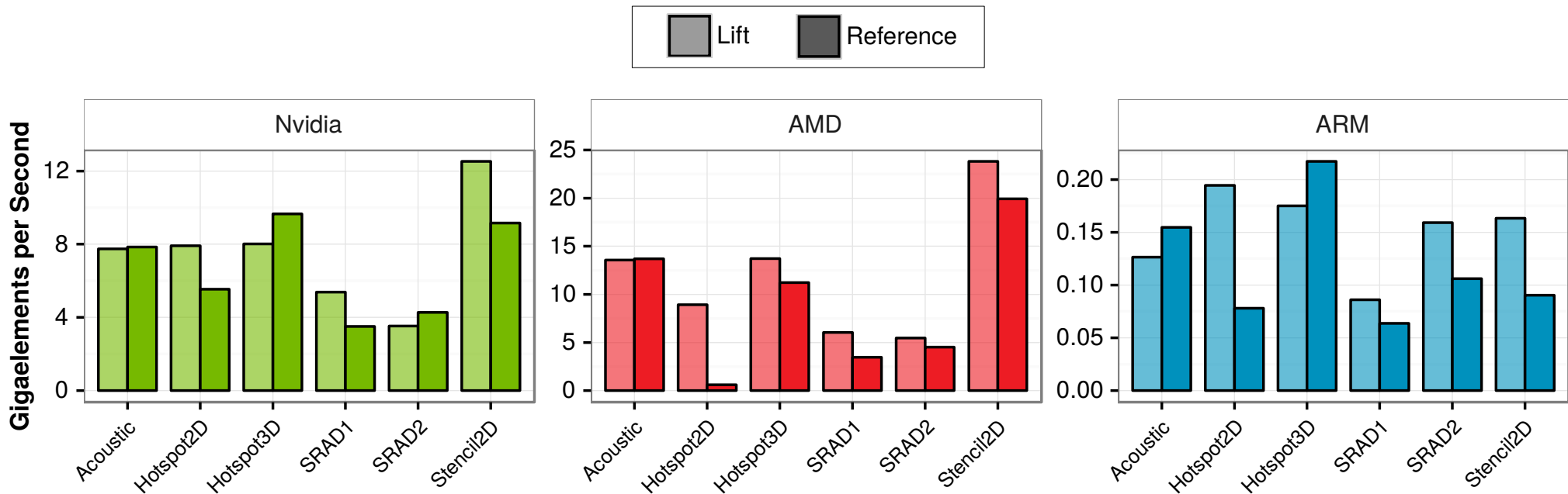
```
stencil = slide(step, size) >> map(f)
```



- ▶ can work in 2D too
- ▶ leverage all existing rewrites

Stencil Computation Results

- ▶ No conceptual changes in the compiler
 - just two new primitives (slide, pad)
 - one extra rewrite rules



Summary

- Rewrite rules define a search space
 - formalisation of algorithmic and optimisation choices
- High performance
 - on par with highly-tuned code

Future Directions

- How to search efficiently the space
 - machine-learning
- Look at other applications
 - e.g. neural network (convolution, recurrent)
- Support for different target programming models
 - MPI, FPGAs, ...

Purpose of the Tutorial

- ▶ Convince you to take up some of the ideas with you
 - functional IR
 - rewriting for optimisation
- ▶ Show you what Lift can do
 - OpenCL code generation
 - Rewriting exploration
- ▶ Find new potential collaboration
 - FPGA backend
 - MPI work
 - new application domain



www.lift-project.org



Michel Steuwer
Lectuer
Glasgow Uni.



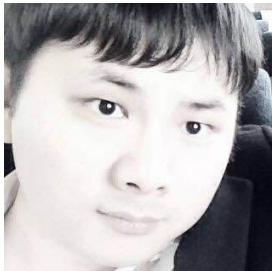
Toomas Remmelg
PhD student
Edinburgh Uni.



Adam Harris
PhD student
Edinburgh Uni.



Bastian Hagedorn
PhD student
Muenster Uni



Lu Li
Postdoc
Edinburgh Uni.



Federico Pizzuti
PhD student
Edinburgh Uni.



Larisa Stolzfus
PhD student
Edinburgh Uni.



Naums Mogers
PhD student
Edinburgh Uni.

supported by:
Oracle Labs
Huawei Innovation Research Program FLAGSHIP
Google Faculty Research Award
Microsoft Research