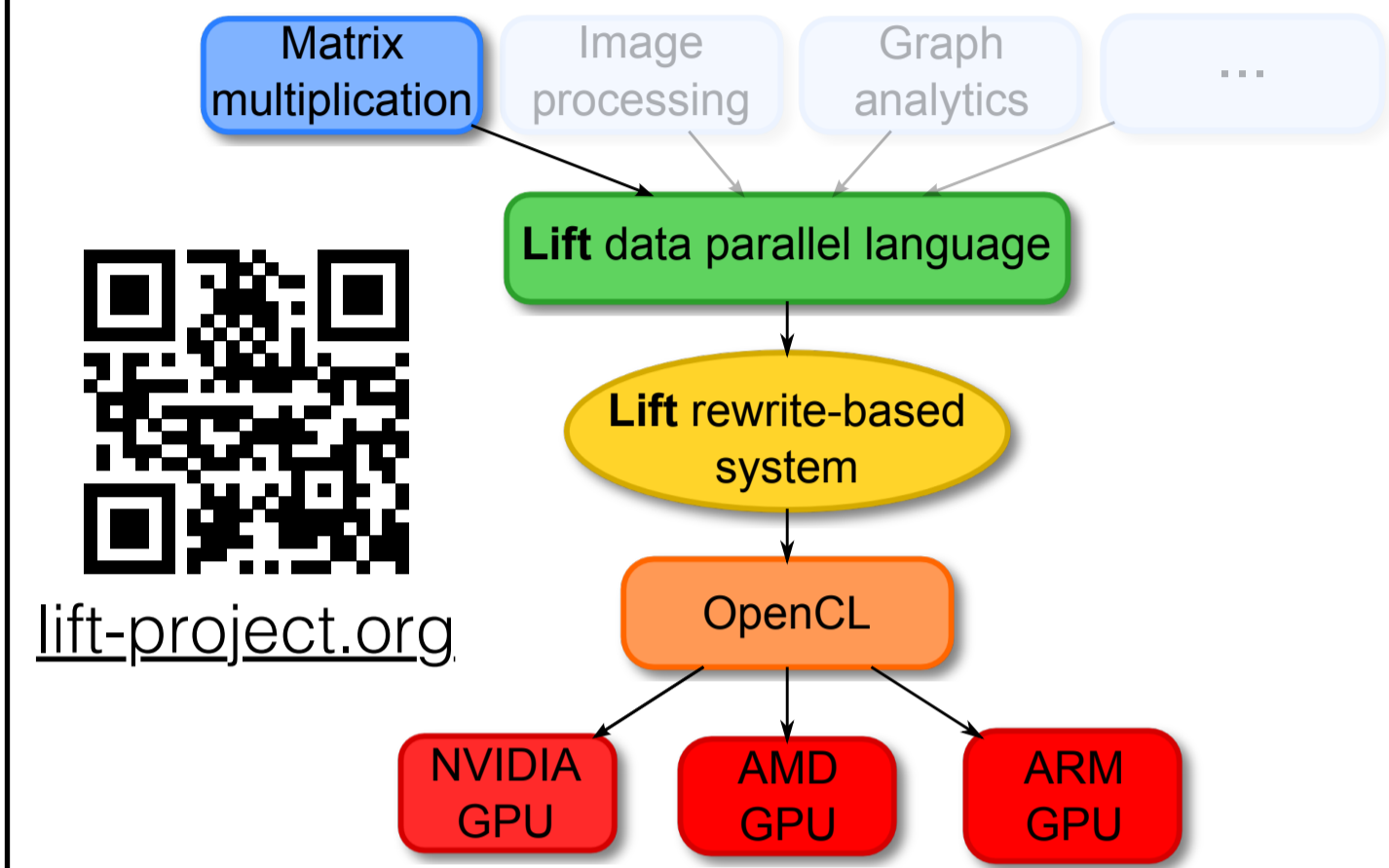


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Lift is a novel approach to achieve high performance on parallel accelerators



Starting from a high-level language, rewrite rules are used to derive optimised implementations.

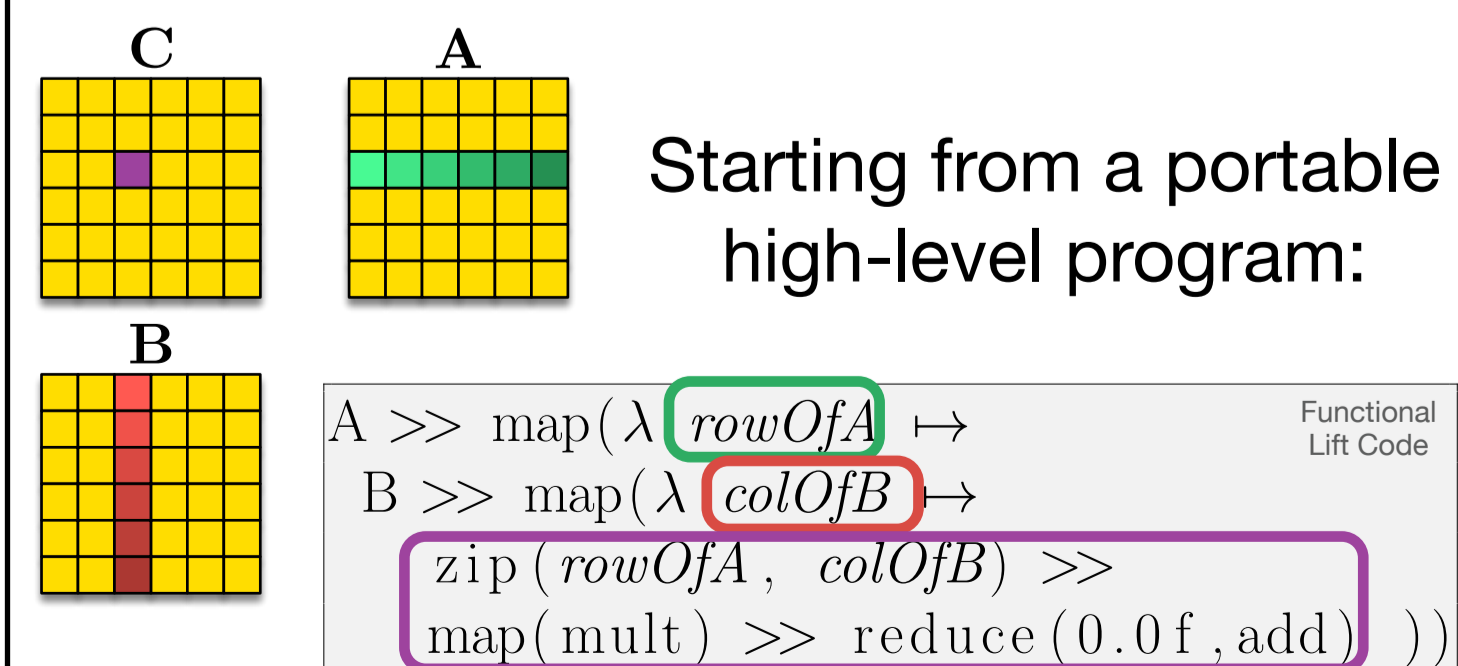
Goal: Achieving Performance Portability, i.e. high performance across different GPUs

Inherent Limitation of Auto-Tuning

Chooses between a fixed number of optimisations and tuning parameters

⇒ Falls short on new architectures!
Performance portability cannot be achieved by only using auto-tuning

Matrix Multiplication in the Lift Data Parallel Language



Lift Rewrite-Based System: Compilation and Optimisation Using Provably Correct Rewrite Rules

Algorithmic rewrite rules express optimisation choices

Split-join rule
 $map(f) \Rightarrow split(k) \gg map \ map(f) \gg join$

$A \gg split(m) \gg map(\lambda rowsOfA \mapsto rowsOfA \gg map(\lambda rowOfA \mapsto B \gg map(\lambda colOfB \mapsto zip(rowOfA, colOfB) \gg map(mult) \gg reduce(0.0 f, add))) \gg join$

```

1 for (int i = 0; i < M/2; i++) {
2   for (int l = 0; l < 2; l++) {
3     for (int j = 0; j < N; j++) {
4       for (int k = 0; k < K; k++) {
5         temp[k + 2*K*N*i + K*N*l + K*j] =
6           mult(A[k + K*l + 2*K*i], B[k + K*j]);
7       }
8     }
9   }
10 }

```

Map-map interchange rule

$X \gg map(\lambda x \mapsto Y \gg map(\lambda y \mapsto f)) \Rightarrow Y \gg map(\lambda y \mapsto X \gg map(\lambda x \mapsto f)) \gg transpose$

$A \gg split(m) \gg map(\lambda rowsOfA \mapsto B \gg map(\lambda colOfB \mapsto rowsOfA \gg map(\lambda rowOfA \mapsto zip(rowOfA, colOfB) \gg map(mult) \gg reduce(0.0 f, add))) \gg transpose \gg join$

```

1 for (int i = 0; i < M/2; i++) {
2   for (int j = 0; j < N; j++) {
3     for (int l = 0; l < 2; l++) {
4       for (int k = 0; k < K; k++) {
5         temp[k + 2*K*N*i + K*N*l + K*j] =
6           mult(A[k + K*l + 2*K*i], B[k + K*j]);
7       }
8     }
9   }
10 }

```

More algorithmic rules

OpenCL specific rules express mapping choices

OpenCL specific rules

$map(f) \Rightarrow mapGlb_{\{0,1,2\}}(f)$
 $map(f) \Rightarrow mapSeq(f)$
 $f \Rightarrow toGlobal(f)$
 $f \Rightarrow toPrivate(f)$
 $map(f) \Rightarrow asVector(n, b) \gg map(vectorize(n, f)) \gg asScalar$

```

A >> split(m) >> mapGlb_{0,1,2}(λ nRowsOfA ↦
B >> split(n) >> mapGlb_{0,1,2}(λ mColsOfB ↦
zip(transpose(nRowsOfA) >> split(k),
transpose(mColsOfB) >> split(k)) >>
reduceSeq(inIt = make2DArray(n,m, 0.0 f) >>
toPrivate(mapSeq(mapSeq(id))))
λ (accTile, (tileOfA, tileOfB)) ↦
zip(accTile, transpose(tileOfA)) >>
mapSeq(λ (accRow, rowOfTileOfA) ↦
mapSeq(λ (acc, colOfTileOfB) ↦
zip(rowOfTileOfA >> asVector(k)
colOfTileOfB >> asVector(k)) >>
mapSeq(dot) >> reduceSeq(acc, add)
) >> join
) >> toGlobal(mapSeq(mapSeq(mapSeq(id))))
>> transpose() >>
map(transpose) >> transpose
) >> join >> transpose
) >> join

```

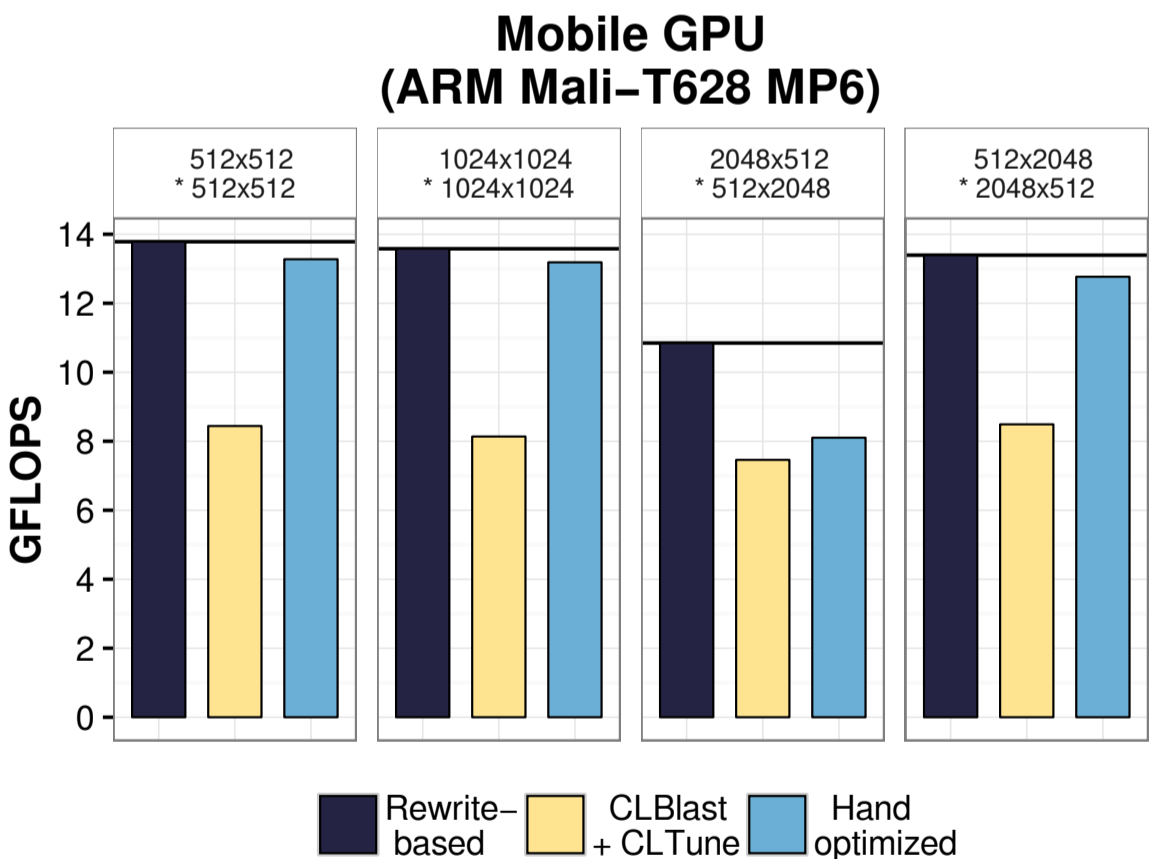
Highly-Optimised Code for the Embedded ARM Mali GPU

```

1 int i = get_global_id(0);
2 int j = get_global_id(1);
3
4 float temp.0; float temp.1;
5 float temp.2; float temp.3;
6 float acc.0; float acc.1;
7 float acc.2; float acc.3;
8
9 for (int k = 0; k < K/4; k++) {
10
11   temp.0 = dot(vload4(k + K*i/2, A),
12             vload4(k + K*j/2, B));
13   acc.0 += temp.0;
14
15   temp.1 = dot(vload4(k + K*i/2, A),
16             vload4(k + K + 2*K*j/4, B));
17   acc.1 += temp.1;
18
19   temp.2 = dot(vload4(k + K + 2*K*i/4, A),
20             vload4(k + K*j/2, B));
21   acc.2 += temp.2;
22
23   temp.3 = dot(vload4(k + K + 2*K*i/4, A),
24             vload4(k + K + 2*K*j/4, B));
25   acc.3 += temp.3;
26 }
27 C[2*N*i + 2*j] = id(acc.0);
28 C[1 + 2*N*i + 2*j] = id(acc.1);
29 C[N + 2*N*i + 2*j] = id(acc.2);
30 C[1 + N + 2*N*i + 2*j] = id(acc.3);

```

High Performance on Mali



Performance Portability Automatic Exploration Using the Lift Rewrite-Based System

